

How functions finish

Direct: `return answer;`

True CPS: `throw k answer;`

uScheme: `(k answer)`

Design Problem: Missing Value

Provide a **witness** to existence:

`(witness p? xs) == x, where (member x xs),
provided (exists? p? xs)`

Problem: What if there exists no such x ?

Solution: A New Interface

Success and failure continuations!

Laws:

```
(witness-cps p? xs succ fail) = (succ x)  
; where x is in xs and (p? x)
```

```
(witness-cps p? xs succ fail) = (fail)  
; where (not (exists? p? xs))
```

Refine the laws

`(witness-cps p? xs succ fail) = (succ x)`
; where `x` is in `xs` and `(p? x)`

`(witness-cps p? xs succ fail) = (fail)`
; where `(not (exists? p? xs))`

`(witness-cps p? '() succ fail) = ?`

`(witness-cps p? (cons z zs) succ fail) = ?`
; when `(p? z)`

`(witness-cps p? (cons z zs) succ fail) = ?`
; when `(not (p? z))`

Coding witness with continuations

```
(define witness-cps (p? xs succ fail)
  (if (null? xs)
      (fail)
      (let ([x (car xs)])
        (if (p? x)
            (succ x)
            (witness-cps p? (cdr xs) succ fail))))))
```

“Continuation-Passing Style”

All tail positions are continuations or recursive calls

```
(define witness-cps (p? xs succ fail)
  (if (null? xs)
      (fail)
      (let ([x (car xs)])
        (if (p? x)
            (succ x)
            (witness-cps p? (cdr xs) succ fail))))))
```

Compiles to tight code

Example Use: Instructor Lookup

```
-> (val 2016f ' ((Fisher 105) (Hescott 170) (Chow 116)))  
-> (instructor-info 'Fisher 2016f)  
(Fisher teaches 105)  
-> (instructor-info 'Chow 2016f)  
(Chow teaches 116)  
-> (instructor-info 'Souvaine 2016f)  
(Souvaine is-not-on-the-list)
```

Instructor Lookup: The Code

```
; info has form: '(Fisher 105)
; classes has form: '(info_1 ... info_n)
(define instructor-info (instructor classes)
  (let (
    [s          ; success continuation
      ]
    [f          ; failure continuation
      ]
    (witness-cps pred
                  classes s f)))
```

Instructor Lookup: The Code

```
; info has form: ' (Fisher 105)
; classes has form: ' (info_1 ... info_n)
(define instructor-info (instructor classes)
  (let (
    [s          ; success continuation
      ]
    [f          ; failure continuation
      ]
    (witness-cps (o ((curry =) instructor) car)
                  classes s f))
  ])
```

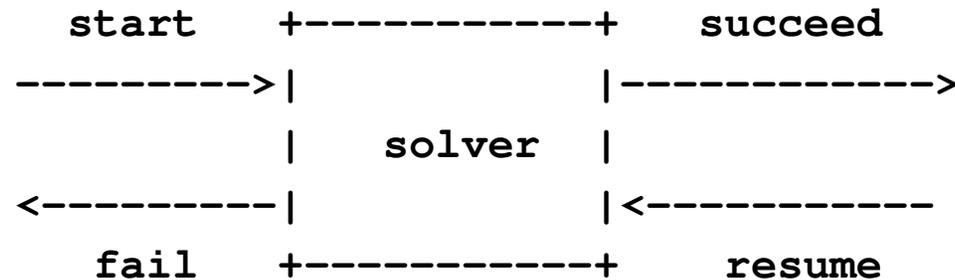
Instructor Lookup: The Code

```
; info has form: '(Fisher 105)
; classes has form: '(info_1 ... info_n)
(define instructor-info (instructor classes)
  (let (
    [s (lambda (info) ; success continuation
          (list3 instructor 'teaches (cadr info)))]
    [f           ; failure continuation
          ])
    (witness-cps (o ((curry =) instructor) car)
                  classes s f)))
```

Instructor Lookup: The Code

```
; info has form: '(Fisher 105)
; classes has form: '(info_1 ... info_n)
(define instructor-info (instructor classes)
  (let (
    [s (lambda (info) ; success continuation
          (list3 instructor 'teaches (cadr info)))]
    [f (lambda ()      ; failure continuation
          (list2 instructor 'is-not-on-the-list))])
    (witness-cps (o ((curry =) instructor) car)
                 classes s f)))
```

Continuations for Search



- start** Gets **partial** solution, **fail**, **succeed**
(On homework, “solution” is assignment)
- fail** Partial solution won't work (no params)
- succeed** Gets improved solution + **resume**
- resume** If improved solution won't work,
try another (no params)

A composable unit!

Continuations for the solver

A big box contains two smaller boxes A and B

There are two ways to wire them up (board)

Imagine A and B as formulas

Imagine A as a formula, B as a *list* of formulas!

Solving a literal

```
(define satisfy-literal-true (x current succ fail)
  (if (bound? x current)
      (if (find x current)
          (succ current fail)
          (fail))
      (succ (bind x #t current) fail)))
```