CS 114: Network Security

Lecture 2

Prof. Daniel Votipka Spring 2023

(some slides courtesy of Prof. Micah Sherr and Prof. Bill Nace)



Administrivia

- Homework 0 due Jan. 26th (Thursday) at 11:59pm
- Homework I, part I due Feb. 2nd at II:59pm
 - Make sure close sockets on CTRL-C
 - There are no additional tests
- If you recently joined the class and do not have access to Gradescope, Piazza, and Box, please email me (<u>daniel.votipka@tufts.edu</u>) or come talk to me after class.

The Plan

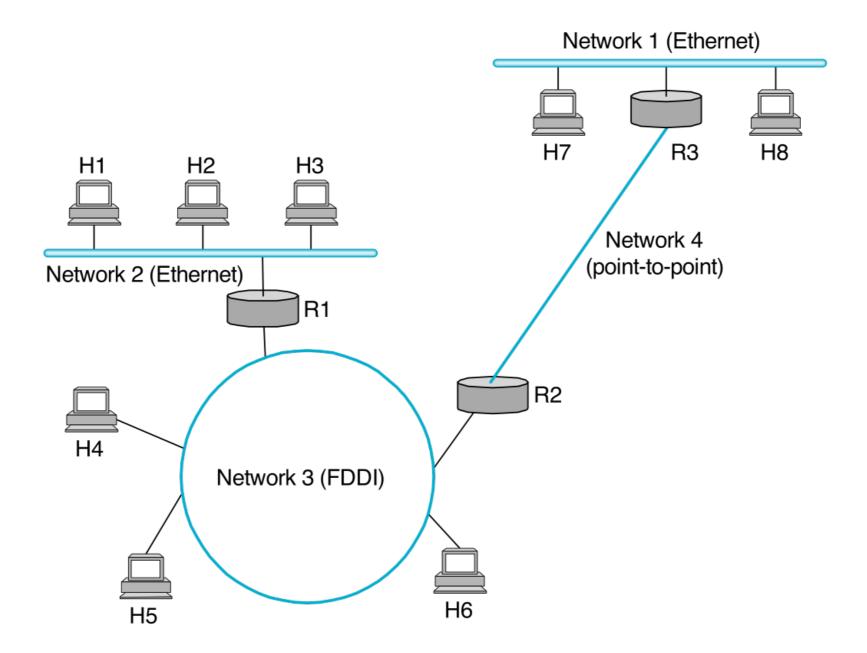
- Introduction to Networking
- The OSI Model
- Network Programming

Intro to Networking

Internet History 101

- DARPA Defense Advanced Research Projects Agency
- ARPANET:
 - World's first operational packet switching network
 - Predecessor of global Internet
 - Started in 1969 with 4 routers @ UCLA, Stanford, UC Santa Barbara, Utah
 - TCP/IP in 1983

Fundamental Goal: An Inter-network

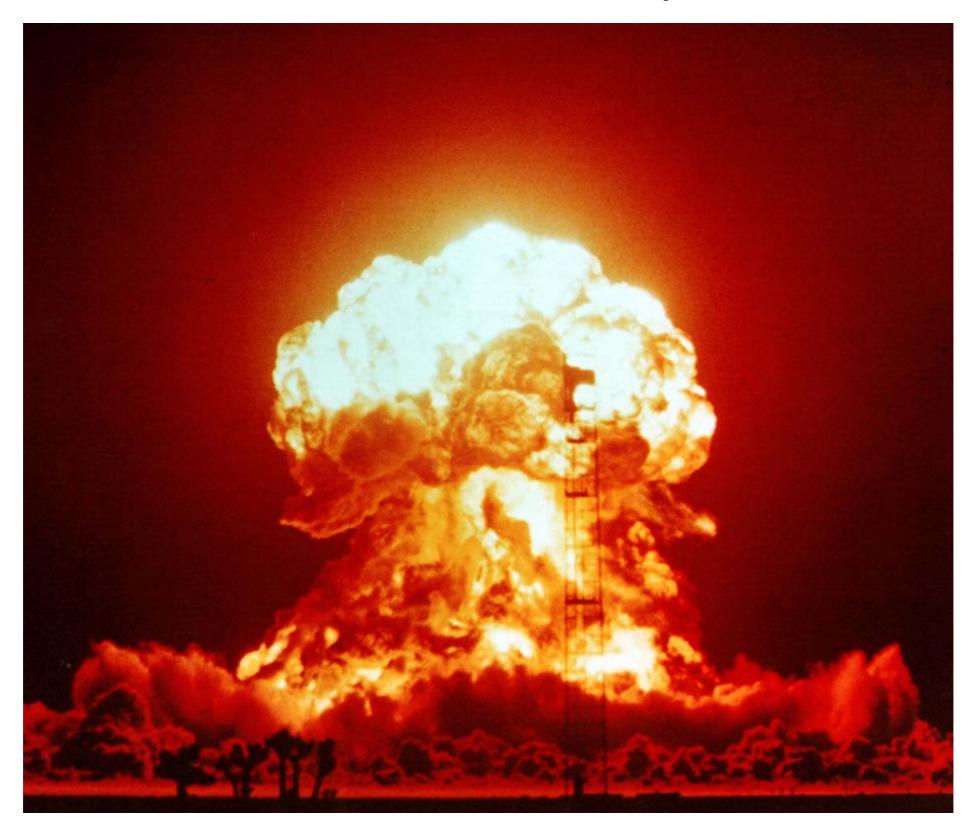


Hn = host, Rn = router/gateway

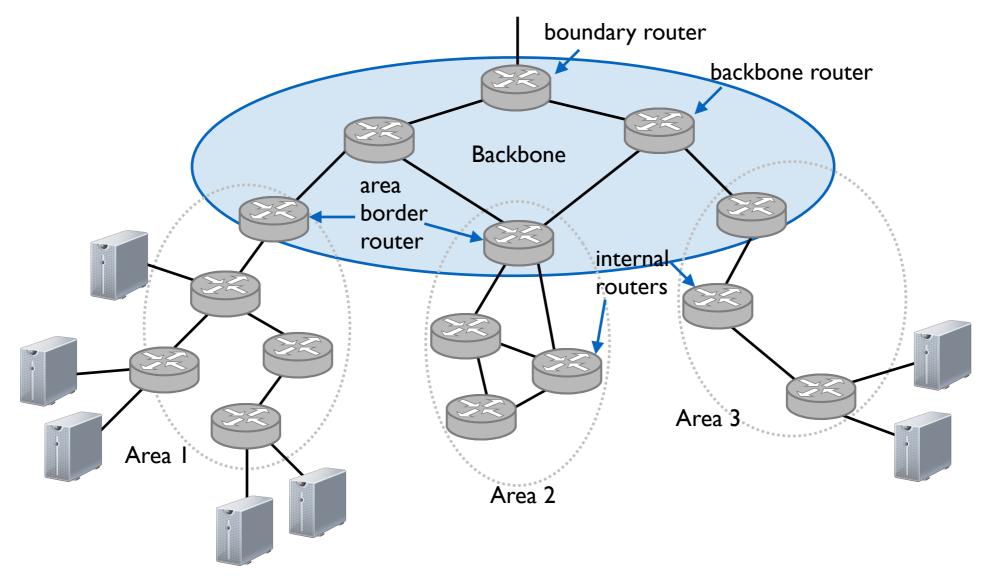
Goals of the Internet

- Fundamental goal: Inter-connect multiple networks of different types (wired and wireless) via store-and-forward gateways
- Second-level goals:
 - Robust in face of failures
 - Support multiple types of services
 - Support a variety of networks
 - Allow distributed management
 - Cost effective

Not a major goal: Route around nuclear explosion

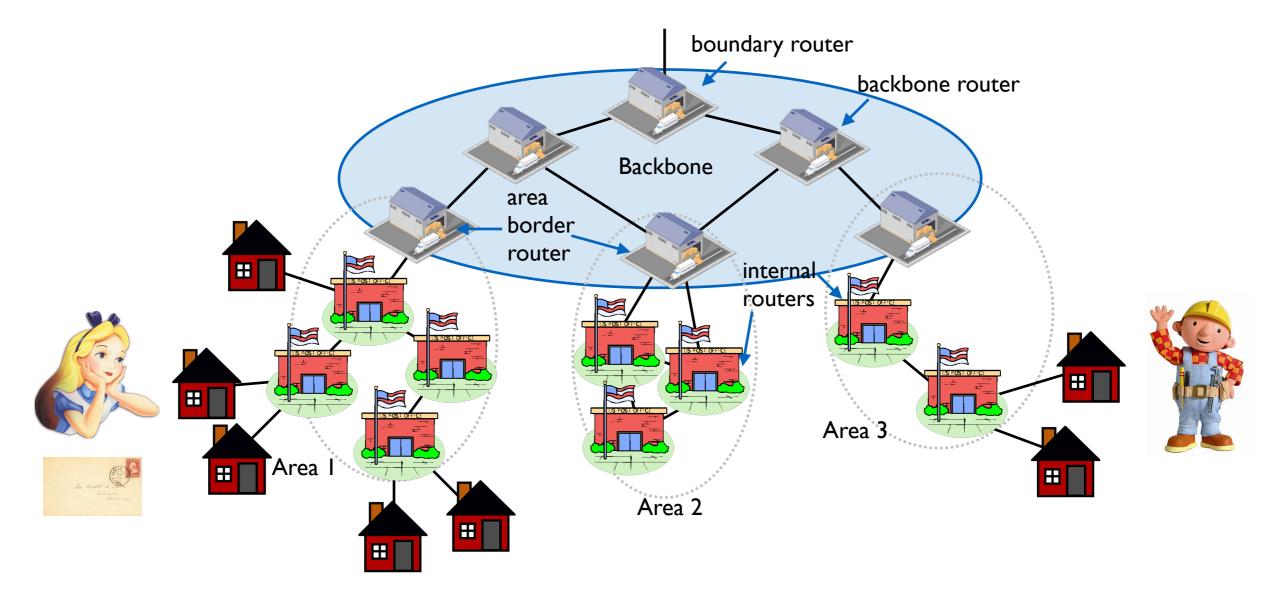


A collection of independently operated autonomous systems (ASes)





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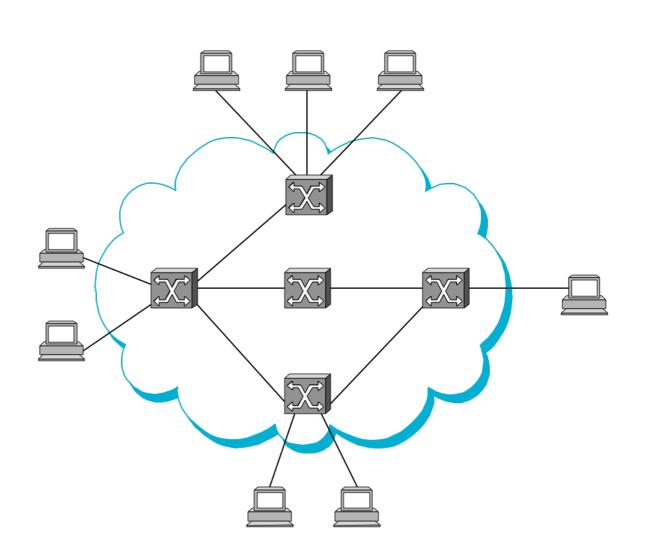


Switched Network

End-hosts connected to switches

• Switches:

- Forwarding nodes
- At least two links
- Also known as bridges or routers



Datagram Packet Switching

- Packets discrete blocks of data
 - Each packet is independently switched
 - Each packet header contains destination address
 - Routing protocol is used to compute next hop
- Example: IP networks
- Advantages:
 - No connection state required
 - Easy to recover from errors
 - Minimal network assumptions

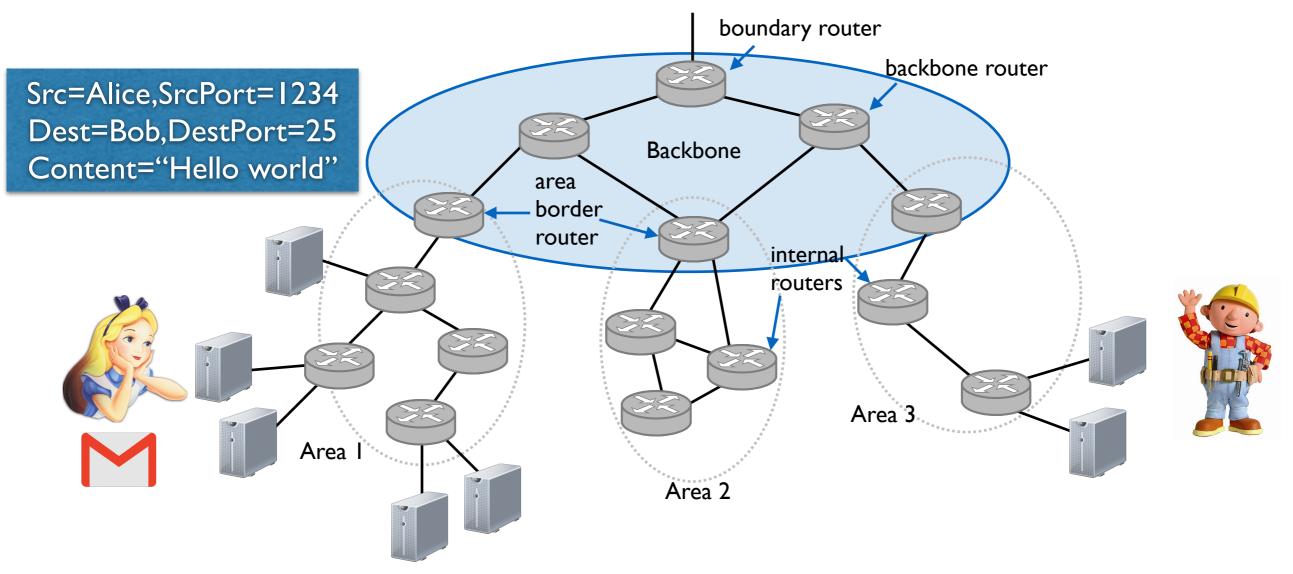
IP Packet Structure

Bits

0 4	4 8	3 1	6 1	9 3	
Version	Length	Type of Service	Total Length		
Identification			Flags	Fragment Offset	
Time to Live Protocol		Protocol	Header Checksum		
Source Address					
Destination Address					
Options					
Data					

https://tools.ietf.org/html/rfc791

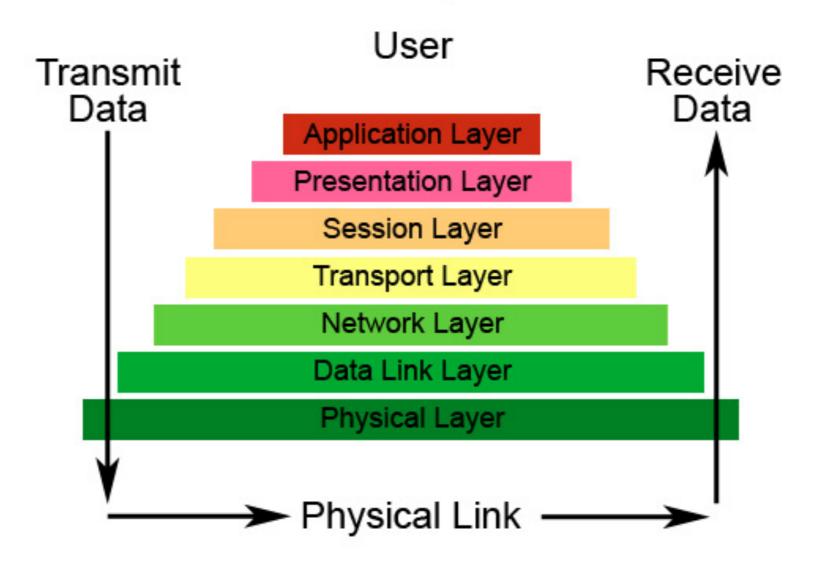
A collection of independently operated autonomous systems (ASes)



The structure of the Internet: The ISO/OSI Architecture

- ISO/OSI architecture describes functions of communication systems
 - ISO: International Standards Organization
 - OSI: Open System Interface
- Breaks communication systems into layers, with each layer serving as an *abstraction* to layer above

The Seven Layers of OSI



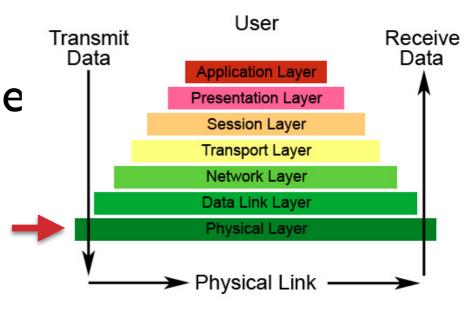
Three Properties of Layers

- Service: what a layer does
- Service interface: how to access the service (interface for layer above)
- **Protocol**: set of rules and formats that govern the communication between two network boxes

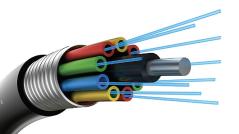


Physical Layer (I)

- Service: move information between two systems connected by a physical link
- Interface: specifies how to send a bit
- Protocol: coding scheme used to represent a bit, voltage levels, duration of a bit
- Examples: coaxial cable, optical fibe links.

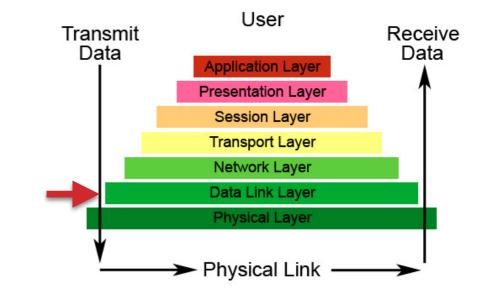


The Seven Layers of OSI



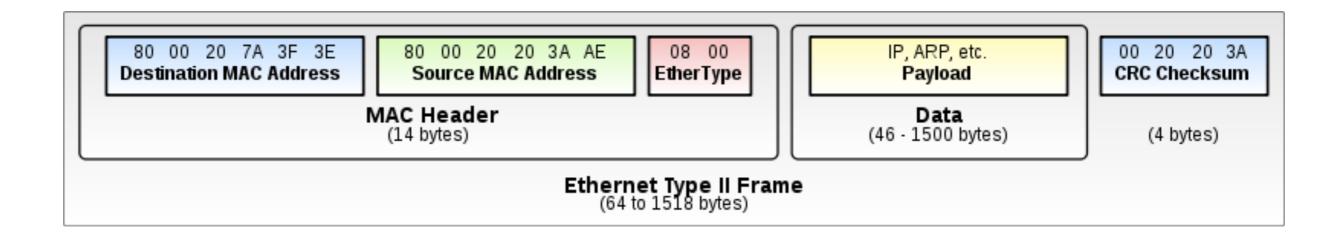
The Seven Layers of OSI

Datalink Layer (2)

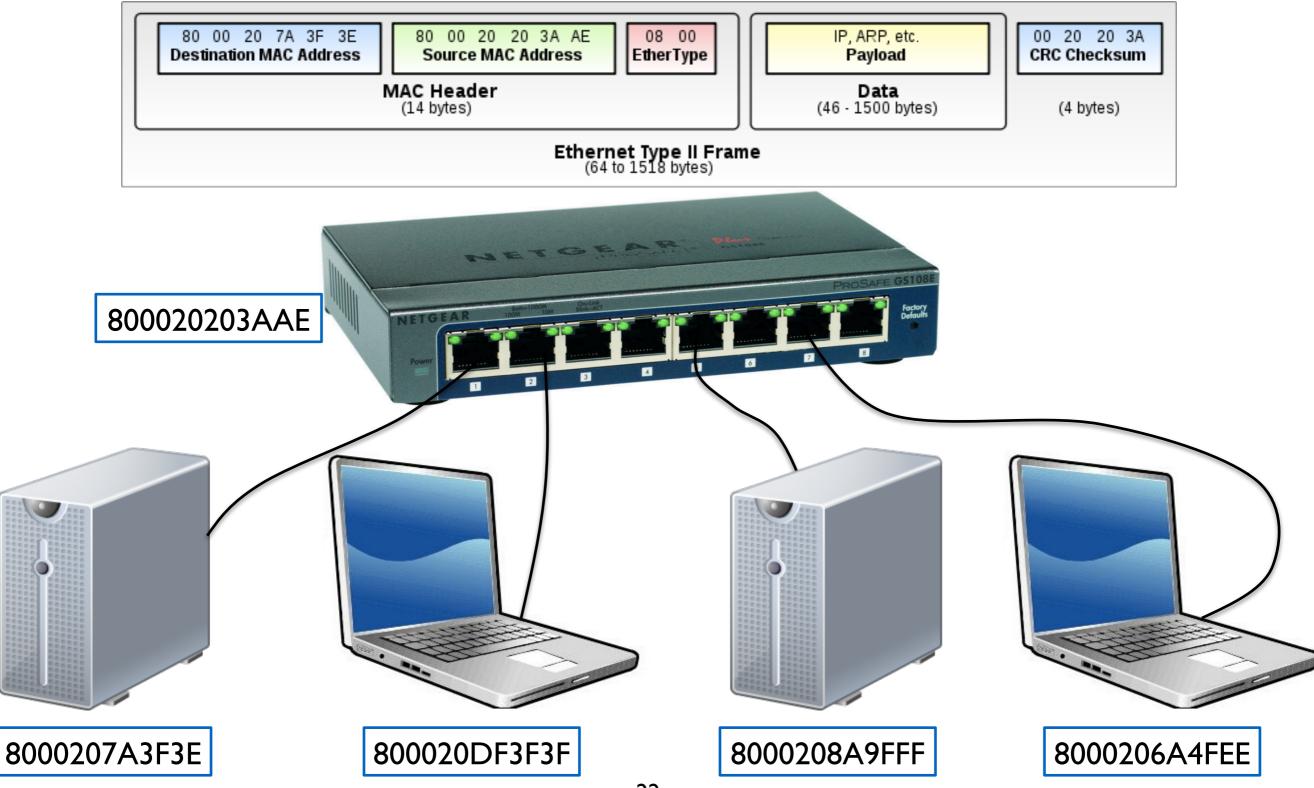


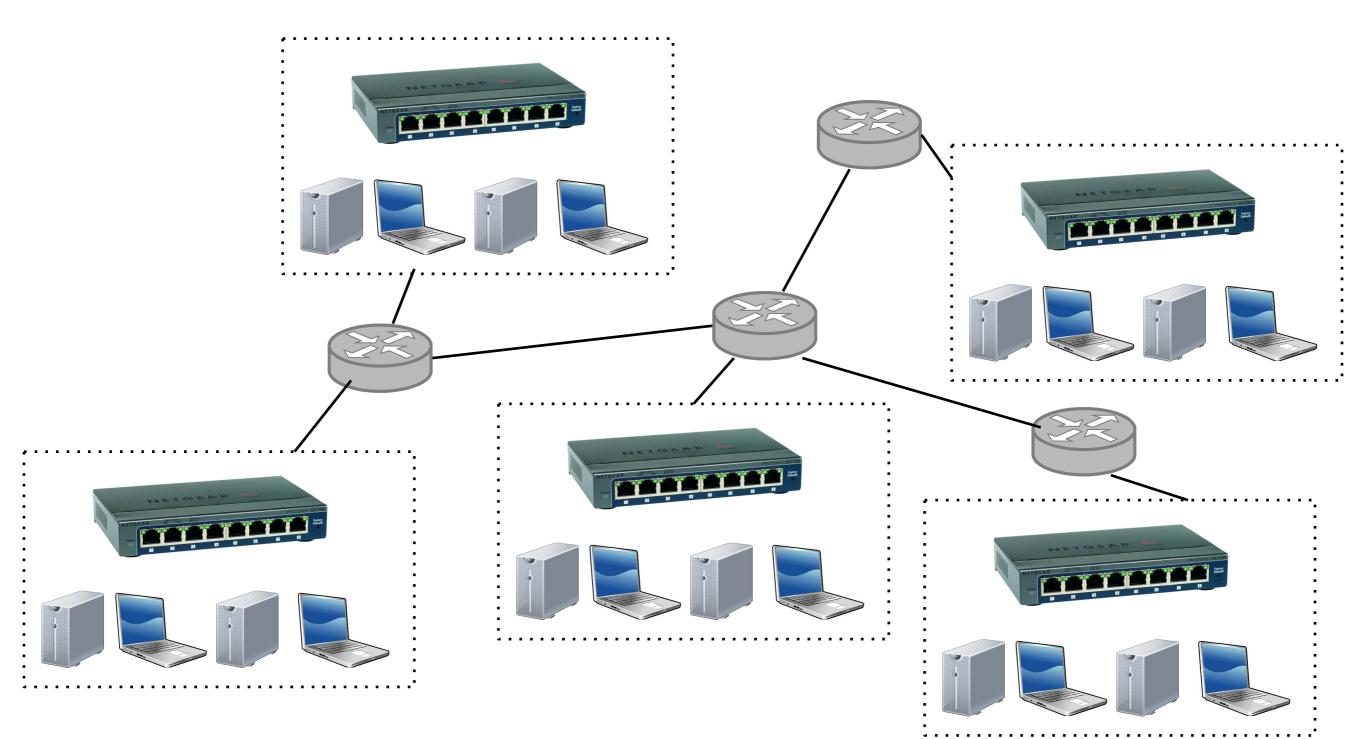
- Service:
 - Combine stream of bits into **frames** (attach frame separators)
 - Send data frames between peers
- Interface: send a data unit (frame) to a machine (MAC address) connected to the same physical media
- Protocol: Medium Access Control (MAC)
 - Examples: Ethernet (LAN), 802.11 (wireless)

Ethernet Frame



An Ethernet Network



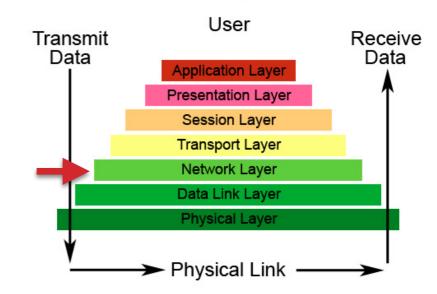


Network Layer (3)

Service:

- Deliver a packet to specified network destination
- Perform segmentation/reassembly
- Interface: send a packet to a specified destination (network address)
- Protocol: define global unique addresses; construct routing tables
 - Example: IP

The Seven Layers of OSI



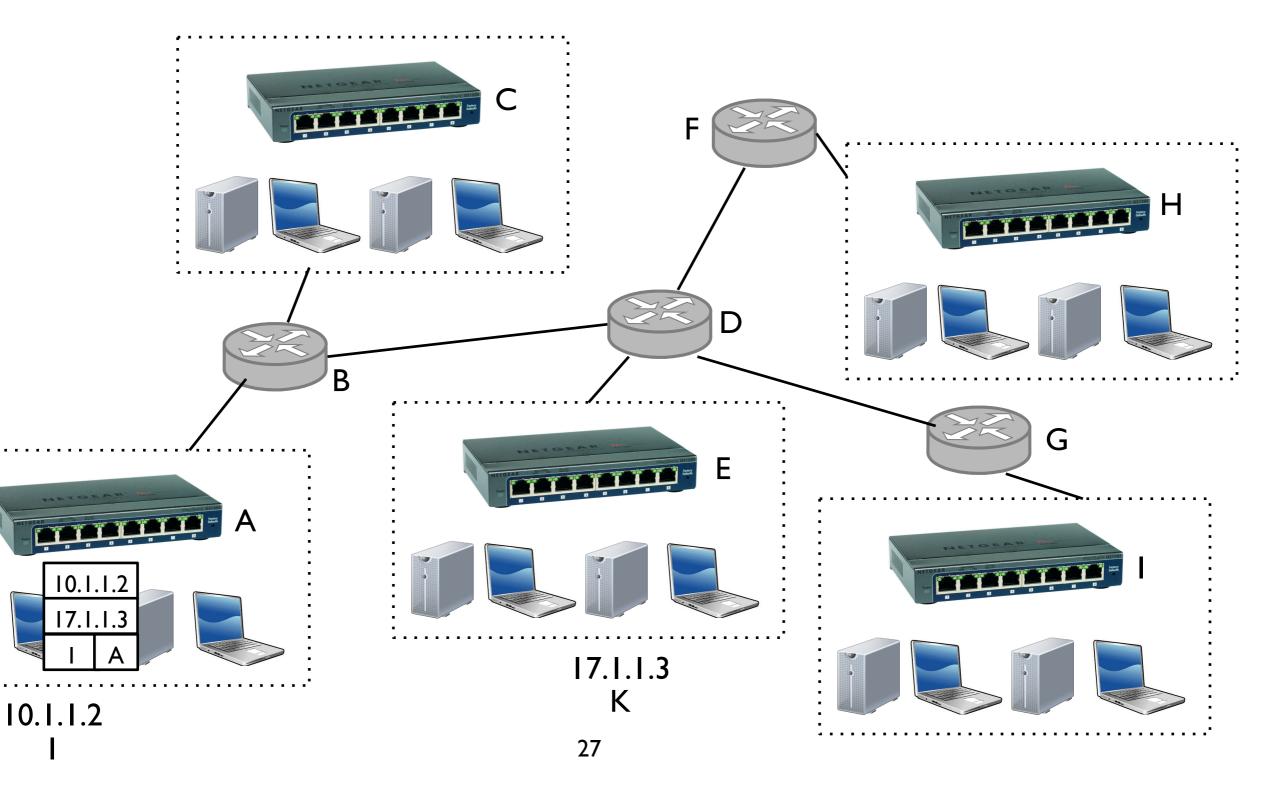
Internet Protocol (IP)

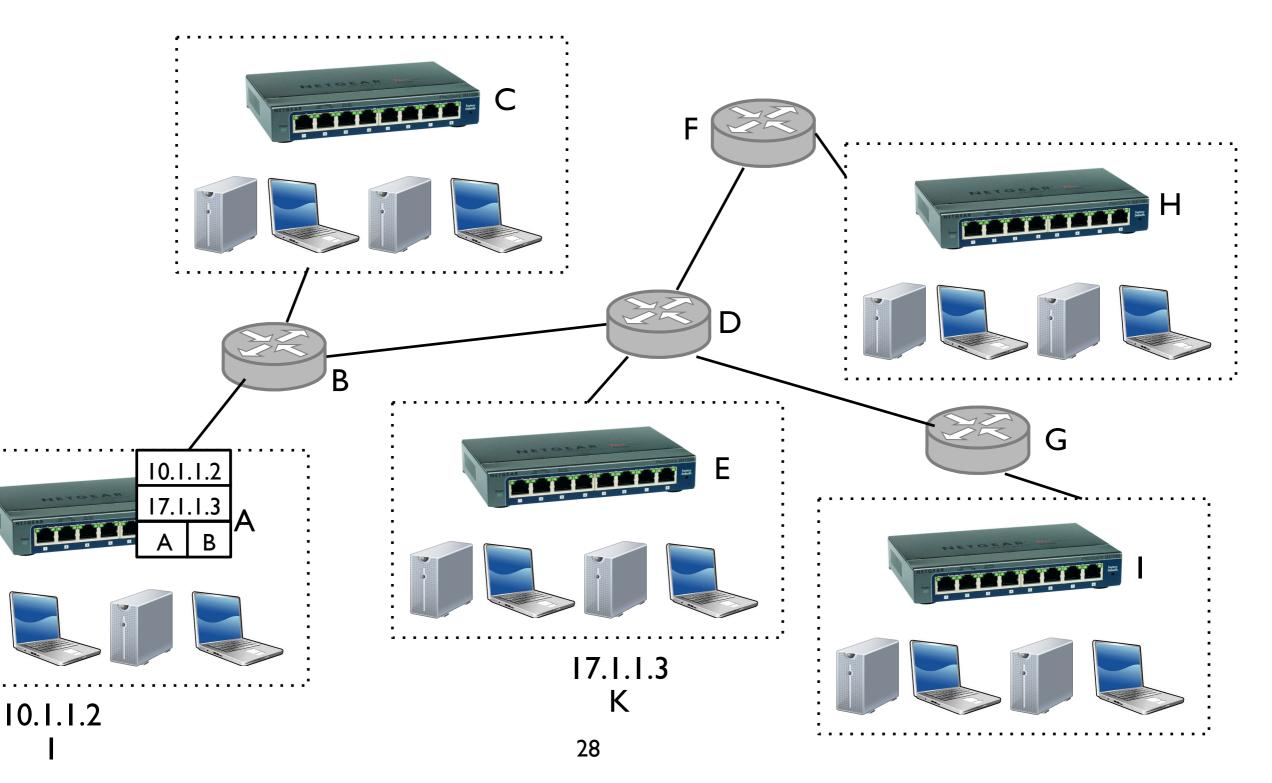
- Allows networks to interoperate
 - Any network technology that supports IP can exchange packets
- Allows applications to function on all networks
 - Applications that can run on IP can use any network

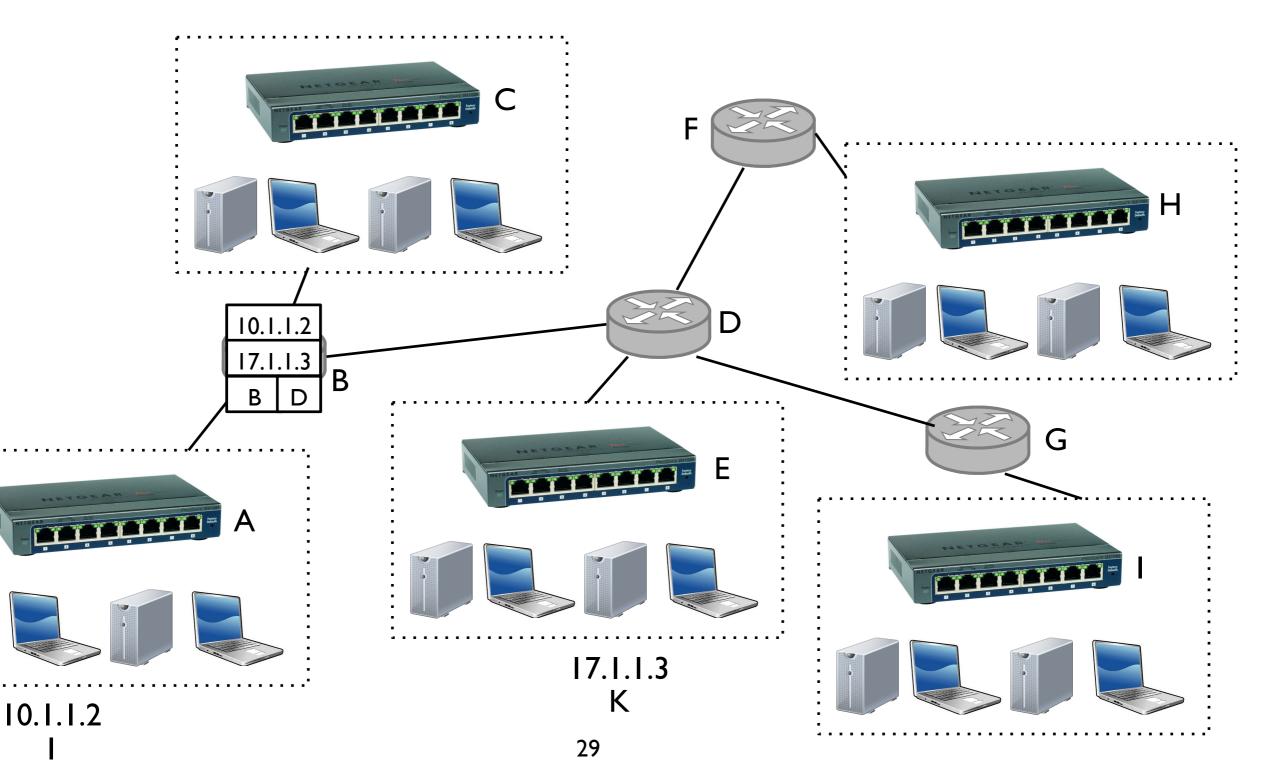
IP

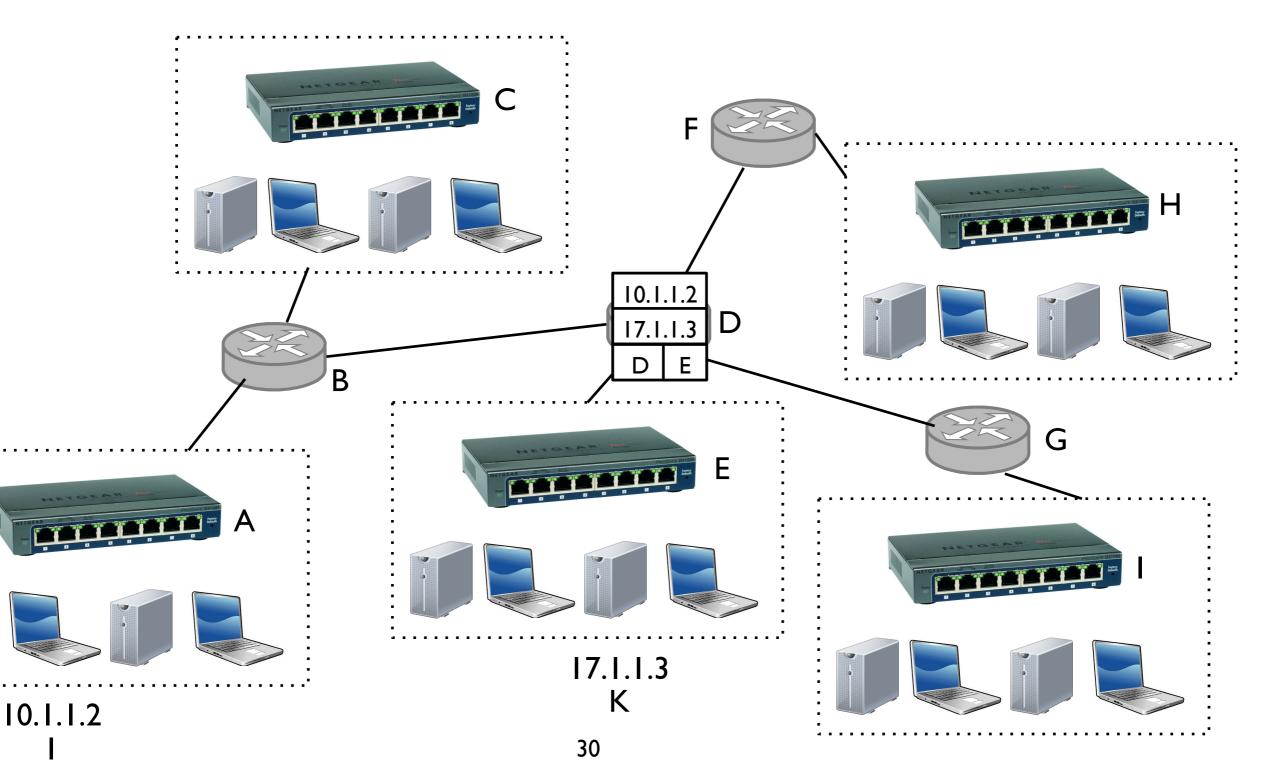
4-bit	8-bit		16-bit	32-bit	
Ver.	Header Length		Type of Service	Total L	ength
Identification			1	Flags	Offset
Time To Live		Protocol		Checksum	
Source Address					
Destination Address					
Options and Padding					

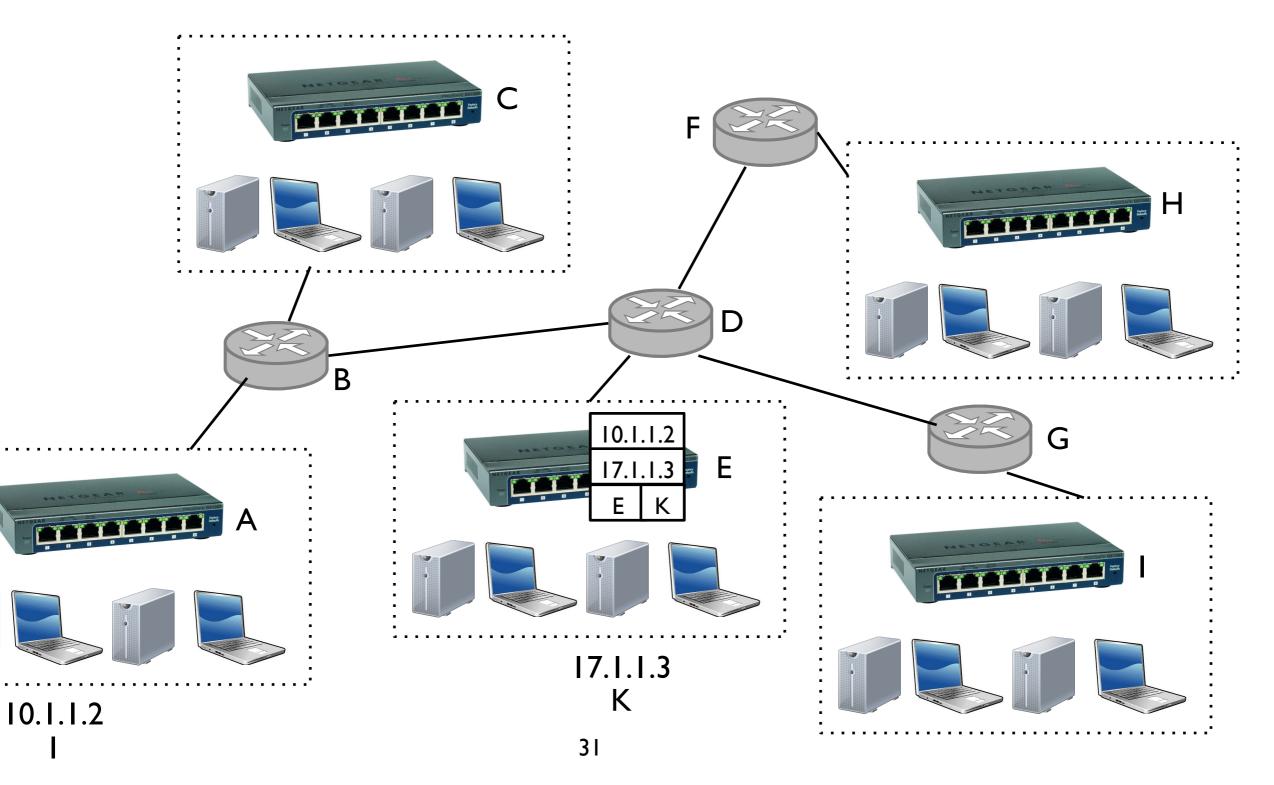
(this is what makes the Internet work)

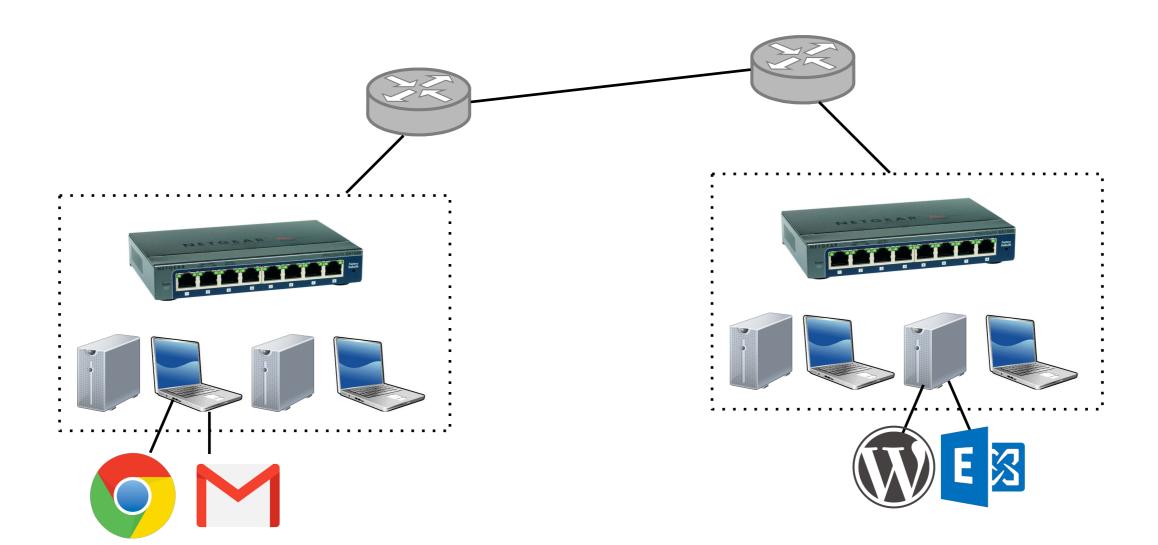








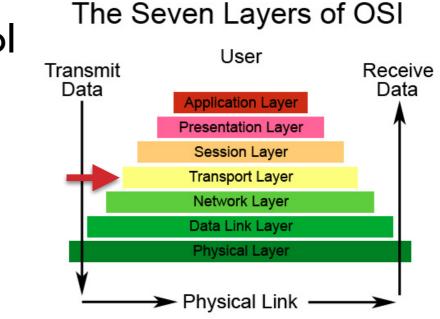




Transport Layer (4)

• Service:

- Process-to-process channels
- Demultiplexing (via ports) to different processes
- Optional: error-free and flow-controlled delivery
- Interface: send message to specific destination (address and port)
- Protocol: implements reliability and flow control
 - Examples:
 - UDP (unreliable)
 - TCP (reliable)



Ports are for multiplexing



Port Number	Service
22	ssh (secure terminal)
25	Email server
80	Web server (HTTP)
443	Secure web server (HTTPS)

Destination port tells OS to which application (service) it should send incoming message

Ports are for multiplexing

- Internet Assigned Numbers Authority (IANA) defines list of reserved ports
- In reality, program can listen to any port in [1,2¹⁶)
 - Ports below 1024 usually require admin privileges
 - For new apps, best to choose a high-numbered port (e.g., 9999 for UnencryptedIM)

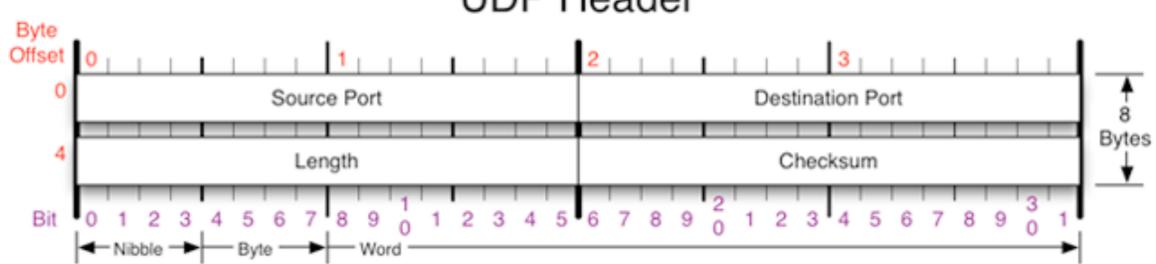
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Destination port tells OS to which application (service) it should send incoming message

User Datagram Protocol (UDP)

- Unreliable transport
- Provides integrity protection
- Doesn't handle:
 - out of order delivery
 - lost packets
 - duplication

When might this be ok?

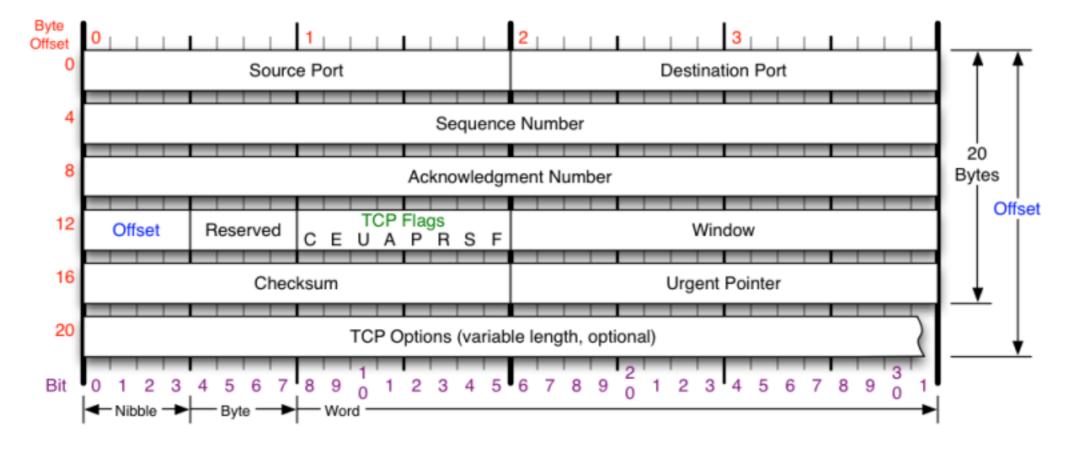


UDP Header

Transmission Control Protocol (TCP)

- Handles loss, duplication, and out-of-order delivery
 - relies on ACKnowledgements
- Provides flow control: prevent congestion
- Provides congestion control: handle congestion

TCP



TCP Flags								
С	Е	U	A	Ρ	R	s	F	
		C	Con	ges	tion	Wi	ndow	
С	0x8	0 F	Red	uce	d (0	CWI	R)	
Е	0x4	0 E	CN	Ec	ho	(EC	È)	
U	0x2	20 L	Irge	ent				
Α	0x1	0 A	ck					
Р	0x0	8 F	usł	n				
R	0x0	4 F	lese	et				
S	0x0	2 S	Syn					

F 0x01 Fin

Congestion Notification

ECN (Explicit Congestion Notification). See RFC 3168 for full details, valid states below.

Packet State	DSB	ECN bits	
Syn	00	11	
Syn-Ack	00	0 1	
Ack	01	0.0	
No Congestion	01	0.0	
No Congestion	10	0.0	
Congestion	11	0.0	
9			
Receiver Response	11	01	
Sender Response	11	11	

TCP Options

0 End of Options List 1 No Operation (NOP, Pad) 2 Maximum segment size 3 Window Scale 4 Selective ACK ok 8 Timestamp

Checksum

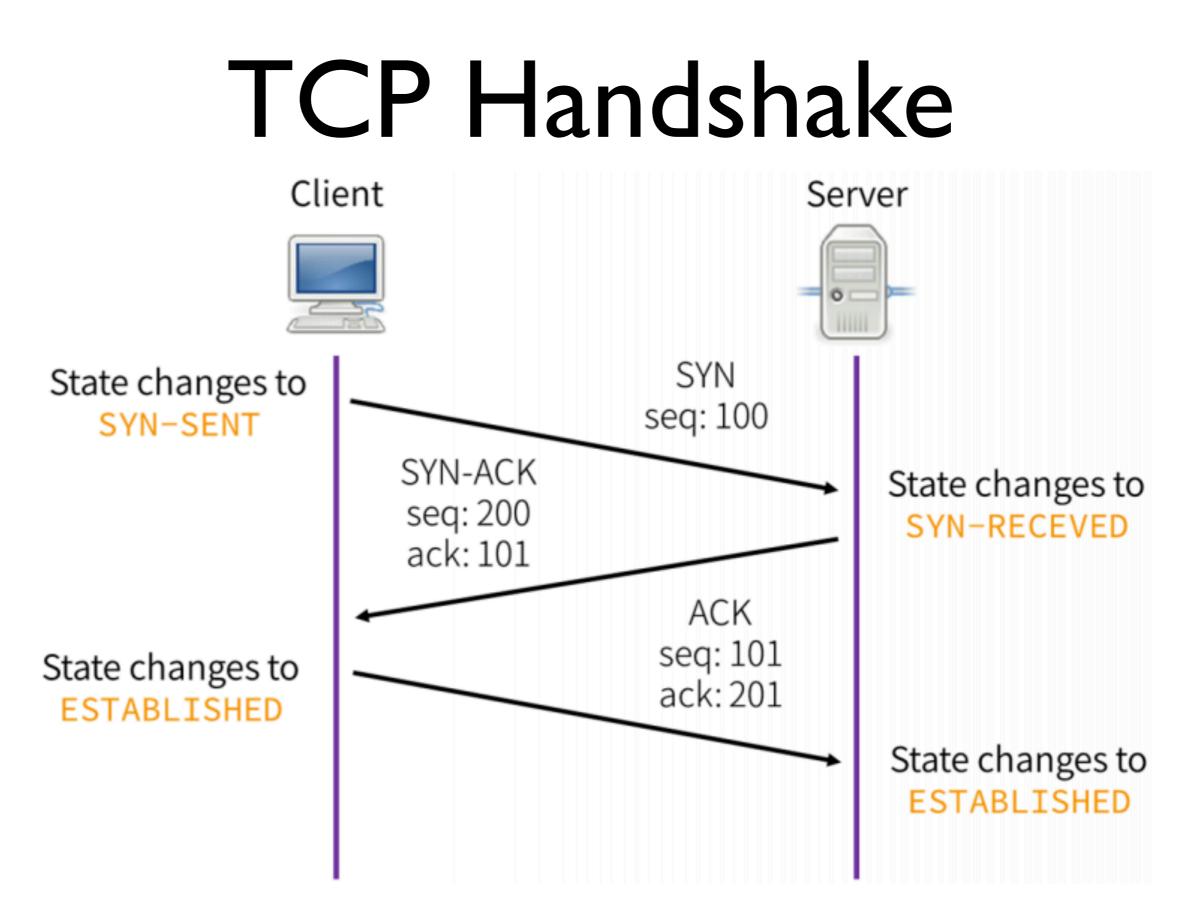
Checksum of entire TCP segment and pseudo header (parts of IP header)

Offset

Number of 32-bit words in TCP header, minimum value of 5. Multiply by 4 to get byte count.

RFC 793

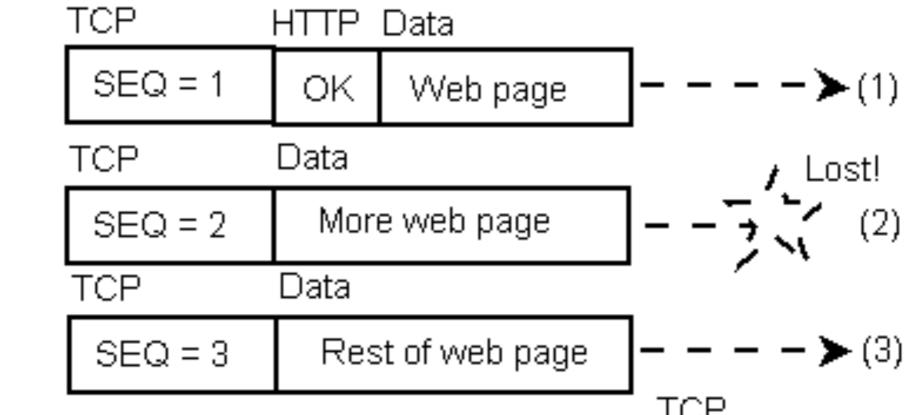
Please refer to RFC 793 for the complete Transmission Control Protocol (TCP) Specification.



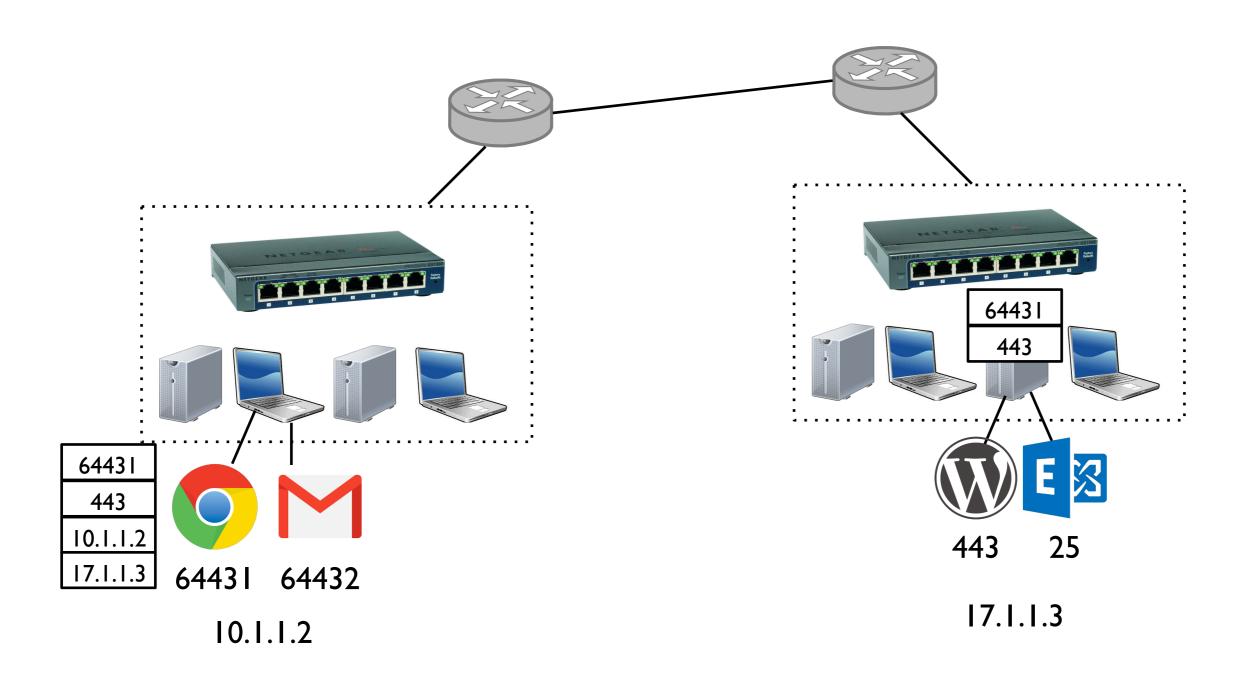
Error Handling

Web server

Web browser



https://reisinge.net/notes/net/basics



Session (5) and Presentation (6) Layers

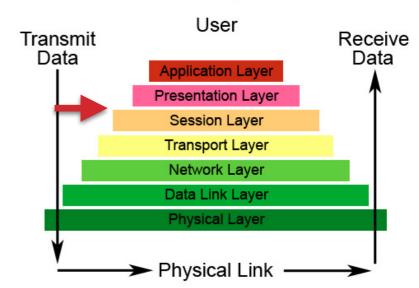
Session service:

- Combines different transport streams for each application
- E.g. audio and video stream in a teleconferencing application

• Presentation service:

 Converts data formats between various representations to provide a standard interface for the application layer **Reality:** These aren't really used.

The Seven Layers of OSI



Application Layer (7)

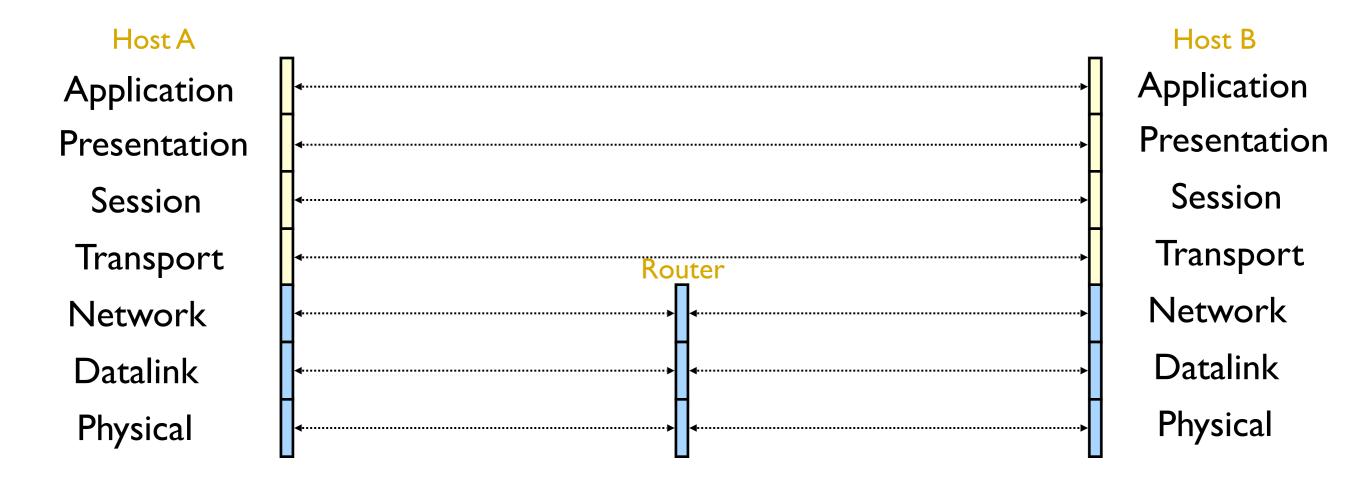
- Service: any service provided to the end user
- Interface: depends on the application
- Protocol: depends on the application
- Examples: FTP, Telnet, SSH, Email, HTTP

Demo!

<u>F</u> ile <u>E</u> dit <u>V</u> iew <u>G</u> o <u>C</u> apture <u>A</u> nalyze <u>S</u> tatistics Telephony <u>T</u> ools <u>I</u> nternals <u>H</u> elp							
	•						
Filter: tcp.stream eq 0 Expression Clear Apply Save							
No. Time Source Destination Protocol Length Info	Â						
1 0.00000000 192.168.2.69 141.161.20.3 TCP 78 50088	>						
7 0.044971000 141.161.20.3 192.168.2.69 TCP 74 http:	> 5						
8 0.044993000 192.168.2.69 141.161.20.3 TCP 66 50088	>						
31 0.471489000 192.168.2.69 141.161.20.3 HTTP 434 GET /:	ma						
40 0.520304000 141.161.20.3 192.168.2.69 TCP 66 http:	• 5 U						
41 0.521241000 141.161.20.3 192.168.2.69 TCP 343 [TCP 3	seg						
42 0.521284000 192.168.2.69 141.161.20.3 TCP 66 50088							
) •						
▷ Uptions: (12 bytes), No-Uperation (NUP), No-Uperation (NUP), Timestamps ▷ [SEQ/ACK analysis] ▼ Hypertext Transfer Protocol	1						
<pre>> GET /images/seal_1.gif HTTP/1.1\r\n</pre>							
Host: www.cs.georgetown.edu\r\n	-17						
Connection: keep-alive\r\n							
Accept: image/webp,*/*;q=0.8\r\n							
User-Agent: Mozilla/5.0 (Macintosh; Intel Mac OS X 10_9_4) AppleWebKit/537.36 (KHTML, like Gecko)	c 🛡						
DNT: 1\r\n	*						
•)+						
0000 e0 3f 49 59 ac 10 b8 8d 12 00 0e 5e 08 00 45 00 .?IY^E.	À						
0010 01 a4 d6 c6 40 00 40 06 fd fb c0 a8 02 45 8d a1@.@E							
0020 14 03 c3 a8 00 50 0a b6 2f a3 24 4d 68 98 80 18P /.\$Mh							
0030 20 10 d6 b2 00 00 01 01 08 0a 41 82 79 68 22 67	¥						
⊖ Ø File: "code/mydump.pcap" 2 Packets: 262 Displayed: 10 Mar Profile: Default							

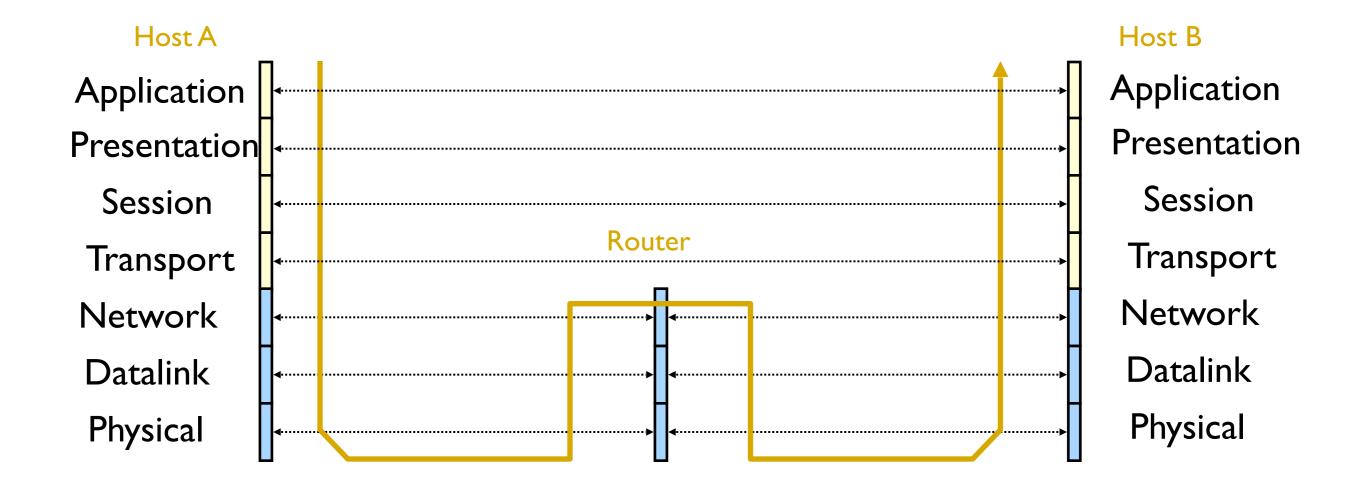
Who Does What?

- Lower three layers are implemented everywhere
- Next four layers are implemented only at hosts



Physical Communication

 Communication goes down to physical network, then to peer, then up to relevant layer



What about security?

- Where is confidentiality, integrity, and authenticity?
- No relevant "security" fields in IP, TCP, or UDP headers.
- Why not?



Network Programming

- The operating system provides an *interface* for sending/ receiving network packets
- A **socket** is a descriptor for network communication
- As a client, you connect your socket to a remote host, and read/write to that socket as you would a file
- As a server, you listen and accept incoming connections, and read/write to that socket as you would a file

Internet communication via sockets

Src=HostA,SrcPort=1234 Dest=HostB,DestPort=1025 Content="Hello world"





HostA

```
import socket
s = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
# connect to HostB on port 1025.
s.connect( ("HostB",1025) )
# Let the Internet worry about how my messages get there
s.send('hello world')
s.close()
```

Internet communication



Src=HostA,SrcPort=1234 Dest=HostB,DestPort=25 Content="Hello world"

Hello world



HostB

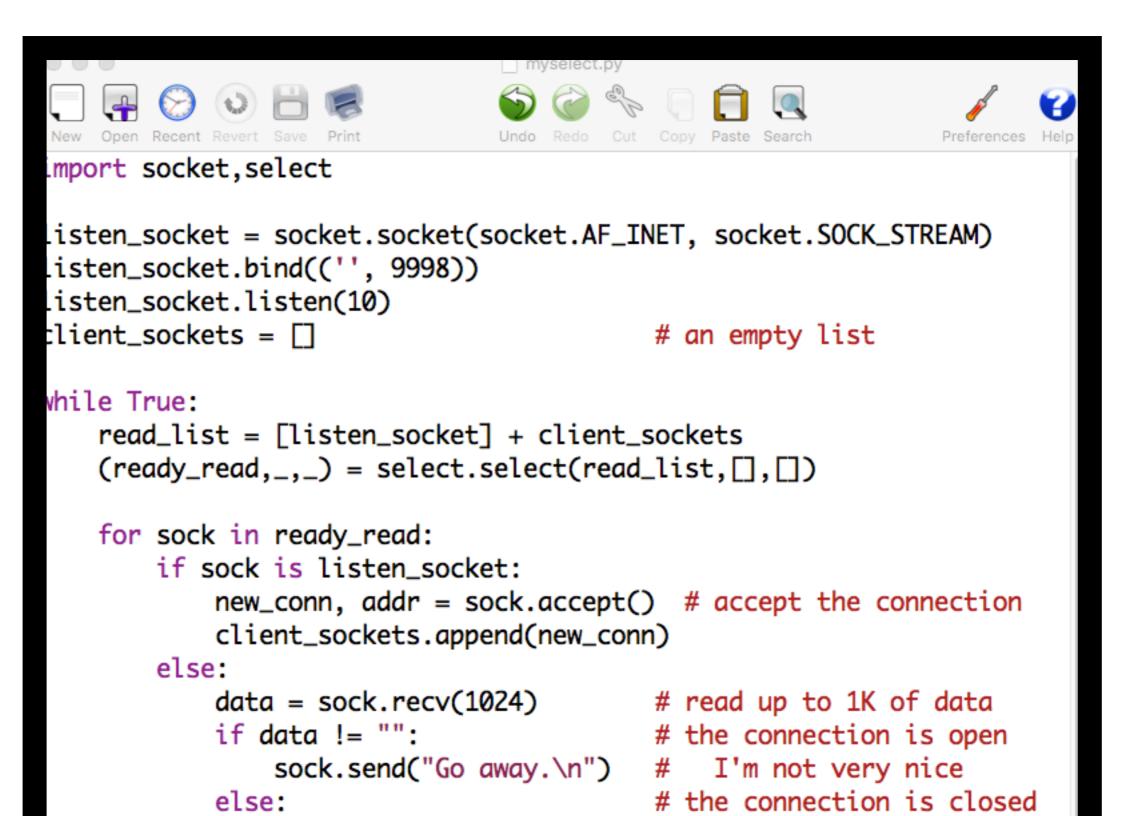
HostA

```
import socket
s = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
s.bind(('', 1025)) # this socket is bound to my port 1025
s.listen(1) # specify the "backlog" for this socket
conn, addr = s.accept() # wait and accept the connection
data = conn.recv(1024) # read the content of the message
print data;
conn.close()
```

Network Programming

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- A **socket** is a descriptor for network communication
- As a client, you connect your socket to a remote host, and read/write to that socket as you would a file
- As a server, you listen and accept incoming connections, and read/write to that socket as you would a file
- read()/recv() is a blocking operation; to wait for input from multiple sources, use select





```
client_sockets.remove(sock)
sock.close()
```

:--- myselect.py All (7,0) (Py Outl)