

# COMP163 Homework Assignment 4

## Due Thursday, November 1, 2018

### Reading:

Read the chapter on Voronoi diagrams/Delaunay triangulations in your text, if you haven't already. Read the comparable sections in the lecture notes. Begin reading in the handouts and in the text about higher dimension convex hulls.

### Problems:

1. The *Gabriel Graph*,  $GG$ , of a set  $S$  in  $E^2$  is defined as follows: let disk  $(p_i, p_j)$  be the circle having  $p_i p_j$  as a diameter; the Gabriel Graph of  $S$  has an edge between  $p_i$  and  $p_j$  in  $S$  if and only if disk  $(p_i, p_j)$  contains no point of  $S$  in its interior. Show that  $p_i$  and  $p_j$  have an edge in  $GG$  if and only if this edge both appears in the Delaunay triangulation of  $S$  and crosses its dual Voronoi edge. Use this to show that the Gabriel graph for any set  $S$  can be constructed in  $O(n \log n)$  time, and also show the optimality of your algorithm. Give an example of a set  $S$  for which the Gabriel Graph and the Delaunay triangulation are NOT the same.
2. In the Euclidean space  $E^d$  of coordinates  $x_1, x_2, \dots, x_d$ , for any real number  $p$  such that  $1 \leq p \leq \infty$ , the  $L_p$ -distance of two points  $q_1$  and  $q_2$  is given by the norm  $d_p(q_1, q_2) = (\sum_{j=1}^d (|x_j(q_1) - x_j(q_2)|)^p)^{1/p}$ . So far, we have been studying the  $L_2$ -distance in  $E^2$  (the plane).

Draw a set of 4-5 points at integer grid points and work out each of the parts below on this small sample. Then generalize if possible.

- (a) In the plane, characterize the Voronoi diagram of a set of  $N$  points in the  $L_1$ -metric.
  - (b) Solve the same problem for the  $L_\infty$ -metric.
  - (c) What is the relationship between the Voronoi diagram in the  $L_1$ -metric and that in the  $L_\infty$ -metric?
3. Given a set  $S$  of  $n$  points in the Euclidean plane and a tolerance  $k$ :
    - (a) describe and analyse an algorithm to report, for each query point  $q$  in the plane, the largest circle centered at  $q$  that contains no point of  $S$  in its interior.
    - (b) describe and analyse an algorithm to report, for each pair of query points  $p, q$ , a path from  $p$  to  $q$  contained entirely within the convex hull of  $S$  such that no point on the path is closer than  $k$  to any point of  $S$ , if such a path exists.
  4. Submit 1-3 rough suggestions of what you might do for a project.