Code Translation and Introduction to Scheme

Judy Stafford Comp 80 – Meeting 4 February 3, 2010

For today, lecture notes available at: http://www.cs.tufts.edu/comp/80/lec

Agenda for the Day

- Administrative Stuff
 - Class list is okay
 - Book ready to float anyone still need it?
 - Moodle...
- ◆ Three Approaches to Code Translation
- Introduction to Scheme
 - a "functional" language

Last time

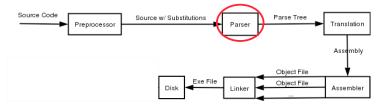
- Syntax
 - Problem: how to precisely describe what a properly formed program looks like
 - Must cover all possible programs
- Solution: formal grammars
 - Inspired by work in linguistics
 - Notation: Backus-Naur Form (BNF)

So, How is this used?

- Abstract syntax tree is the basis of "interpretation" and "compilation"
 - Interpreted languages
 - » Respond to you in real time
 - » You type in statement and get a result
 - ♦ 01
 - » You feed it a list of statements in a file and get a result
 - Compiler
 - » You feed it a file with source code and it converts it to executable code that can be saved and run later

So, again, How is it used?

◆ Compilation (ex: C++) -- Five basic phases



- Interpretation (ex: Scheme)
 - Each line is *parsed* into its smallest operations
 - Each operation is executed
 - Code remains as text you just get the top line of the diagram above
 - » And, you get that every time you run the program!

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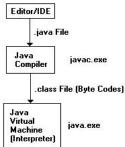
How's Parsing done?

- Requires three phases
 - Lexer read characters in and form tokens
 - » Requires a scanner and a tokenizer
 - » Produces a table of tokens
 - Syntactic analysis
 - » Validates that tokens form a legal expression
 - Semantic analysis
 - » Type checking
 - » Variable def/use relationships
 - » etc.

And then there is Java...

- Java code is "compiled" into byte code
- JVM (Java Virtual Machine) "interprets" the bytecode

"Write once, run anywhere"



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Java Byte Code

```
• Example Java:
```

```
StringBuffer aStringBuffer = new StringBuffer("Hello");
aStringBuffer.append("World !");
```

• Example Byte Code:

```
0  new #8 <Class java.lang.StringBuffer>
3  dup
4  ldc #2 <String "Hello">
6  invokespecial #13 <Method java.lang.StringBuffer(java.lang.String)>
9  astore_110 aload_111 ldc #1 <String "World !">
13 invokevirtual #15 <Method java.lang.StringBuffer append(java.lang.String)>
16 pop
```

- The Byte Code is what is distributed and run
 - When it is run it is interpreted by the JVM for the target platform

Summarizing

- Grammars have been developed to define an unambiguous way for the machine to translate code
- Parse trees are used as an internal representation of the code
- Compilers and interpreters use parse trees to check the code and give feedback to programmer
- Finally the code is turned into executable in one of the three ways

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Intro to Scheme

- ◆ Scheme is a functional language
 - That doesn't mean it has more functions than C++ or Java
 - It means it is more like functions in mathematics

John McCarthy



- ◆ Pioneer in AI
 - Formalize commonsense reasoning
- ◆ Also
 - Proposed timesharing
 - Mathematical theory
- ◆ Lisp
 - stems from interest in symbolic computation (math, logic)

Lisp, 1960

- Look at historical Lisp
 - Example of elegant, minimalist language
 - Not "C" (or Java): a chance to think differently
 - Many general themes in language design
 - Perspective
 - » Some old ideas seem new, but some are just old
- Many different dialects
 - Lisp 1.5, Maclisp, ..., Scheme, ...
 - CommonLisp has many additional features
 - Use "Lisp" to mean family of languages
 - This course: we'll use Scheme

Running programs

- ◆ A modern version of Lisp: Scheme
 - Interactive system: DrScheme
 Download from www.drscheme.org
 - It's available on linux machines
 - We'll do some Scheme programming
- Supports multiple dialects of Scheme
 - Click on "Language" tab then "Choose Language"
 - Choose "Advanced Student"
 - Use "Show Details" to set output style to "quasiquote"
- ◆ I'll demo as we go…

Languages you've seen

Declaration

int x

- Introduces new identifier
- May bind value to identifier, specify type, etc.
- ◆ Expression

(f(x) + 5) / 2

- Syntactic entity that is evaluated
- Has a value, need not change accessible memory
 Change to memory is called a side effect
- Statement

p->next = q

- Imperative command
- Alters the contents of previously-accessible memory

Languages you've seen

A program is...

- Sequence of statements
 - "Do this, then do this, then do this...."
 - Update the state of memory until we have the answer
 - Remember x = x + 1?
- Possibly organized into procedures
 - For convenience
 - For reuse
 - Usually have "parameters" so that they can operate on different inputs

Scheme syntax: atoms

Symbols

foo bar course comp80 Computability 1 2 3 7up

Special symbols

(For boolean truth values)

TRUE FALSE
Historical Lisp: T nil
Scheme: #t #f

DrScheme*: true false

*student language

Scheme syntax: lists

Simple case: list of atoms

```
(1 2 3 4)
(apple orange banana)
(2 apples 3 oranges no bananas)
```

- What do you notice about the contents of lists?
- No distinction between kinds of atoms (no types)
- Lists can contain other lists

```
((2 apples) (3 oranges) (0 bananas))
((the (quick brown) fox) jumped over (the lazy dog))
```

Empty list

() <u>or</u> nil <u>or</u> null

Lists

- Very simple data structure
- BUT, you can build almost anything
- ◆ Examples:
 - A list (duh)
 - A record (or struct)
 - A list of records?
 - A binary tree?

Scheme programs

- Lists and atoms
 - Called "S-expressions"
 - S is for "symbolic"
 - That's it!
- OK, what's the trick?
 - Special way of interpreting lists as code called eval
 - Idea:

```
Given a list (fn a1 a2 a3)
```

Treat fn as a function and a1, a2, and a3 as arguments.

Notice: lists used for both code and data

More details

Function application

$$(fn arg_1 ... arg_n)$$

- First, evaluate each of the arguments (arguments might also be function applications)
- Second, pass list of argument values to function
- Notice: *prefix* notation function name comes first (as opposed to *infix* notation: x + 3)
- Many built-in functions
- Later: make your own functions

Special forms: quote

- Problem:
 - What happens here: (sum (1 2 3 4))
 - What if I don't want to interpret the list as code?
- Quoting a value

```
(quote A) → A also written: 'A
```

- Does not evaluate A
- Just returns A as a value
- Examples:

```
(quote 3) \rightarrow 3
(quote (+ 3 2)) \rightarrow (+ 3 2)
'(+ 3 2) \rightarrow (+ 3 2)
'(To be or not to be) \rightarrow (To be or not to be)
```

List manipulation

- ◆ Two main functions (built-in)
- car also called first
 - Returns the first element of the list
- cdr also called rest
 - Returns everything except the first element
 - Note: the result is always a list
- Use combinations to obtain parts

More lists

- How do I build a list?
- One option: use quote special form
 - What's the problem?
 - Can't compute the list parts:
 '(1 2 (first (3 4 5 6)) does not return '(1 2 3)
- Solution: cons function makes a new list
 - Takes an element and a list
 - Makes a new list with that element on the front
 - (cons 1 '(2 3)) returns (1 2 3)
- ◆ Examples...

Other useful functions

- Tests (predicates)
 - eq?, equal? Atom/List Equality
 (eq? 3 4) → false
 - atom?, null? Is an atom or null? (atom? 'foo) → true
- Logical operators
 - and, or, not Logical operations(not true) → false
- How do we use these tests?

Interesting notes

- No assignment
 - Do we need it?
- Lists are never actually modified
 - We always get a new list
 - Sounds a little crazy, but it's very useful
- No loops
 - No for, while, goto, etc.
 - Problem for next time...
- Notice:
 - If we can build and manipulate lists
 - And lists are used to represent code...

Next time

- More on Scheme and intro to names and types
 - Read Chapters 4 and 5
- Short Scheme assignment due next Thursday
 - Will be posted on moodle asap
 - Will include instructions for downloading and running DrScheme