



COMP80

Programming Languages



Introduction to ML

February 24, 2009

Prelude


- How many cells are in the human body?
 - Approximately 10 to 100 trillion cells
 - Tricky: varies from person to person
 - And: are we counting just *human* cells?
 - Body hosts many kinds of microbes
 - Like what?
 - *Bacteria on skin, intestinal flora*
- Of all the cells, what percentage are human?
 - Approximately 1 in 10 cells!
 - Human cell is 100X to 1000X bigger
- What if we killed all the bacteria in our bodies?
 - We would be dead in about 2 weeks

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Last time


- Types
 - Basic types
 - Boolean, integer, real (float)
 - Composite types
 - Enum, array, records/struct, list, pointer
- Types in languages
 - Static vs dynamic typing
 - Strong vs weak typing
- Today: programming with types



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Shapes problem


- Programming problem
 - Shape: square, circle, triangle
 - Circle – radius
 - Rectangle – width, height
 - Triangle – base, height
 - Given a list of shapes, compute total area
- Let's code it together...



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Scheme issues

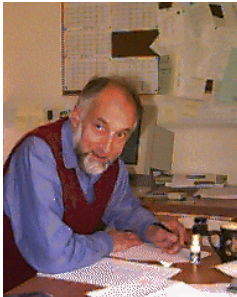

- Data structures
 - Lists of lists of lists
 - How do I know I have a circle, square, rectangle?
 - How do I enforce the structure?
 - How do we check for errors?
 - When do we find out about errors?
- New language called ML
 - Statically typed
 - We can define new types, like "Circle"
 - ML compiler will make sure that functions that expect a circle are always given a circle



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History of ML

- Robin Milner
- Logic for Computable Functions (LCF)
 - Stanford 1970-71
 - Edinburgh 1972-1995
- Meta-Language of the LCF system
 - Theorem proving
 - Type system
 - Higher-order functions

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ML

- Language profile
 - Functional language, like Scheme
 - Strongly typed, like Scheme
 - But:
 - Statically typed
 - Everything has a well-defined type
 - Must pass exactly the right thing into a function

```
(define first-atom (lambda (x)
  (if (list? x) (first-atom (first x))
      x)))
```

- Too restrictive?



Why a new language?

- Discuss types (duh)
 - General issues of type
 - Type checking
 - Provides a framework for understanding other languages
- Advanced topics
 - Type inference
 - Algebraic datatypes
 - Polymorphism
- Very powerful programming techniques



ML basic types

- Primitive types
 - **int**
 - **real**
 - **bool**
 - **string**
- Types are unambiguous
 - Lexicographic form identifies type
 - Examples:
 - 3 : **int**
 - 3.0 : **real**
 - "hello, world" : **string**

Type syntax:
<expression> : <type>



ML syntax

- Scheme
 - Extremely uniform
(function arg1 arg2 ... argN)
- ML
 - Infix notation for many operations
 - 3 + 2 instead of (+ 3 2)
 - More familiar function call notation:
myfun arg1 arg2 ... argN
 - What does that mean?
 - No need for quote – e.g., '(1 2 3)
 - Few, if any, parentheses – only when needed
 - Downside?
 - Have to worry about operator precedence



ML functions

- Scheme:
 - (define double (lambda (x) (* 2 x)))
 - Or shorthand:
(define (double x) (* 2 x))
- ML
 - fun <name> <arg1> ... <argN> = <expression>;
- Examples:
 - fun double x = 2 * x;
 - fun coins q d n p = (q*.25) + (d*.10) + (n*.05) + p;
 - fun doubleadd x y = double x + double y;

Note precedence: function call
higher than addition



Lists

- Similar to Scheme
 - Different syntax
[1, 2, 3, 4]
 - What else do you think is different about ML lists?
 - What is the type of [1, 2, 3, 4]?
[1, 2, 3, 4] : int list
 - ML lists must contain elements of the same type
 - What if they didn't?
 - We could not check at compile-time that computations on the list (e.g., adding up the elements) were legal
- Examples...



List operations

- Scheme
 - How do we build a list in Scheme?
 - Cons
- ML has cons
 - Infix operator `::`
 - Scheme: `(cons 1 '(2 3 4))` → `'(1 2 3 4)`
 - ML: `1::[2, 3, 4]` → `[1, 2, 3, 4]`
 - Scheme empty list: `'()`
 - ML empty list: `nil`
 - Scheme: `'(1 2 3) = (cons 1 (cons 2 (cons 3 '())))`
 - ML: `[1, 2, 3] = 1::2::3::nil`



Type checking

- Who has made this error:
 - Given `'(1 2 3)` and `'(4 5 6)`
 - Want to make `'(1 2 3 4 5 6)`
 - Using `(cons '(1 2 3) '(4 5 6))`
 - Get `'((1 2 3) 4 5 6)`
- In ML
 - What is `[1, 2, 3]::[4, 5, 6]`?
 - Not allowed!
 - Result would be `[[1, 2, 3], 4, 5, 6]`
 - What is the type?



Tuples

- Group of values of non-uniform type
 - Like a record or struct
 - Syntax
 - `(1, 3.141, "foobar")`
 - What is the type of that value?
 - Called the cross product: `int × real × string`
 - Conceptually:
 - *A tuple of this type is a member of the set of all possible triples that consists of one int, one real, and one string.*
 - **Note:** a specific number of items
 - **Note:** order matters!



Combinations

- Combinations? You bet...
- Lists of lists
 - `[[1, 2, 3], [4], [5, 6]]`
 - Type?
- List of tuples
 - `[("Guyer", 100), ("Ramsey", 85), ("Hescott", 60)]`
 - Type?
- Tuple of lists
 - `("Guyer", [75, 85, 60, 100])`
 - Type?



Value declarations

- Bind a value to an identifier
 - **Note:** *not a "variable"*
 - Form: `val <name> = <expression>;`
- Examples
 - `val mylist = [1, 2, 3, 4];`
 - `val mygrade = ("Guyer", 100);`
 - `val grades = [mygrade, ("Ramsey", 85), ("Hescott", 60)];`
- Local bindings
 - Similar to Scheme let
 - `let val x = 2 * 3 in x*4;`



Lists and tuples

- How do we get stuff out of lists and tuples?
 - You might think: functions (like first and rest)
 - You'd be wrong
 - **Question:** what if the left-hand side of a val declaration was not just a simple variable?
 - `val first::rest = [1,2,3,4]`
- **Idea:** *pattern matching*
 - General form: `val <pattern> = <expression>;`
 - Two parts:
 - Match structure of <pattern> with the value of <expression>
 - Bind the names in the pattern to the parts
 - Example: `first = 1, rest = [2,3,4]`
 - Also called *unification*



Pattern matching

- Examples:

Breaking open a tuple:

- `val (name, grade) = ("Guyer", 100);`
- `val mygrade = ("Guyer", 100);`
- `val (name, grade) = mygrade;`

Get the first two elements of a list?

- `val thelist = [1,2,3,4];`
- `val first::second::rest = thelist;`

What happens here?

- `val first::rest = mygrade;`
- Clearly, this won't work – what is the problem?
- Problem: types don't match



More ML functions

- Scheme

```
(define (fact x)
  (if (eq? x 1) 1
      (* x (fact (- x 1)))))
```

- ML

- Uses pattern matching on function cases
- General form:

```
fun <pattern> = <expression>
  | <pattern> = <expression> ...etc...
```

```
fun fact 1 = 1
  | fact x = x * fact (x - 1);
```

Order matters:
ML chooses the
first pattern that
matches



More patterns

- Processing lists

```
(define (somefunc L)
  (if (empty? L) base-case
      (combine (first L) (somefunc (rest L)))))
```

- In ML

- What are the two cases?
- Empty and non-empty list
- Patterns?
- nil for base case, first::rest for inductive case

```
fun somefunc nil = basecase
  | somefunc first::rest = combine first rest;
```



More functions on lists

- Length of a list

```
fun len nil = 0
  | len x::xs = 1 + len xs;
```

- Append lists

```
fun append nil ys = ys
  | append x::xs ys = x :: append xs ys;
```

- Reverse a list

```
fun reverse nil = nil
  | reverse (x::xs) = append ((reverse xs), [x]);
why not | reverse (x::xs) = (reverse xs):: [x]; ?
```

- Questions

- How efficient is reverse?



From our Scheme lecture

- Define addup

(Adds up the elements of a list)

- Define double

(Doubles each element of a list)

- Higher-order functions?

- You bet
- **foldr** and **map**



Next time

- More ML types

We'll solve the shapes problem

- Type inference

- Polymorphism

