



COMP 181

Lecture 15

Code generation II



October 26, 2006

Prelude

- Who are these people?
 - Pierre and Pam Omidyar
- And we care because....?
 - Pierre Omidyar is the founder of eBay
 - Donated \$100 million to Tufts
- But there are strings...
 - The money must be invested in **microfinance**
- What is microfinance?
 - Small business loans, typically given to individuals in developing countries
 - Average: \$600 (as low as \$40)


The person who invented microfinance just got the Nobel prize in economics

2

Conference report


- Architectural Support for Programming Languages and Operating Systems
 - Mix of three areas
- Highlights:
 - Spatial scheduling for tiled architectures
 - GPU programming
 - Automatic generation of peephole superoptimizers
 - Combinatorial sketching for finite programs



3

Code generation

- Overview
 - Walk AST, generate code for each construct
 - Nested constructs:
 - Call generate on children as necessary
 - Use registers to pass values up tree
- Key:** order is important
 - Emitted code goes into central list
 - Bottom-up post-order traversal

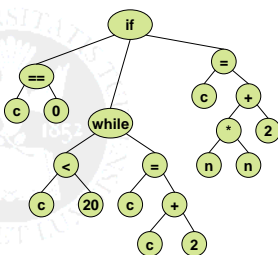



4

Example

```

if (c == 0) {
  while (c < 20) {
    c = c + 2;
  }
} else
  c = n * n + 2;
    
```

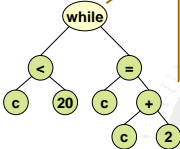



5

Example


```

E = new_label() = L0
T = new_label() = L1
emit( T: )
t = generate(expr)
emit( ifnot_goto t, E )
generate(statement)
emit( goto T )
emit( E: )
    
```



Code

L1:



6

Example

```

E = new_label() = L0
T = new_label() = L1
emit( T: )
t = generate(expr)
emit( ifnot_goto t, E )
generate(statement)
emit( goto T )
emit( E: )
  
```

Code

L1:

7

Example

```

E = new_label() = L0
T = new_label() = L1
emit( T: )
t = generate(expr)
emit( ifnot_goto t, E )
generate(statement)
emit( goto T )
emit( E: )
  
```

Code

L1:

```

t1 = generate(expr1)
t2 = generate(expr2)
r = new_temp()
emit( r = t1 op t2 )
return r
  
```

8

Example

```

E = new_label() = L0
T = new_label() = L1
emit( T: )
t = generate(expr)
emit( ifnot_goto t, E )
generate(statement)
emit( goto T )
emit( E: )
  
```

Code

L1:
R0 = load c

```

t1 = generate(expr1)
t2 = generate(expr2)
r = new_temp()
emit( r = t1 op t2 )
return r
  
```

```

r = new_temp() = R0
emit( r = load v )
return r
  
```

9

Example

```

E = new_label() = L0
T = new_label() = L1
emit( T: )
t = generate(expr)
emit( ifnot_goto t, E )
generate(statement)
emit( goto T )
emit( E: )
  
```

Code

L1:
R0 = load c

```

t1 = generate(expr1)=R0
t2 = generate(expr2)
r = new_temp()
emit( r = t1 op t2 )
return r
  
```

10

Example

```

E = new_label() = L0
T = new_label() = L1
emit( T: )
t = generate(expr)
emit( ifnot_goto t, E )
generate(statement)
emit( goto T )
emit( E: )
  
```

Code

L1:
R0 = load c
R1 = 20

```

t1 = generate(expr1)=R0
t2 = generate(expr2)
r = new_temp()
emit( r = t1 op t2 )
return r
  
```

```

r = new_temp() = R1
emit( r = 20 )
return r
  
```

11

Example

```

E = new_label() = L0
T = new_label() = L1
emit( T: )
t = generate(expr)
emit( ifnot_goto t, E )
generate(statement)
emit( goto T )
emit( E: )
  
```

Code

L1:
R0 = load c
R1 = 20
R2 = R0 < R1

```

t1 = generate(expr1)=R0
t2 = generate(expr2)=R1
r = new_temp() = R2
emit( r = t1 op t2 )
return r
  
```

12

Example

```

E = new_label() = L0
T = new_label() = L1
emit( T: )
t = generate(expr)=R2
emit( ifnot_goto t, E )
generate(statement)
emit( goto T )
emit( E: )

```

```

Code
L1:
R0 = load c
R1 = 20
R2 = R0 < R1
not_goto R2,L0

```

Tufts University Computer Science

Example

```

E = new_label() = L0
T = new_label() = L1
emit( T: )
t = generate(expr)=R2
emit( ifnot_goto t, E )
generate(statement)
emit( goto T )
emit( E: )

```

```

Code
L1:
R0 = load c
R1 = 20
R2 = R0 < R1
not_goto R2,L0

```

```

r = generate(expr2)
l = lgenerate(expr2)
emit( store *l = r )
return r

```

Tufts University Computer Science

Example

```

E = new_label() = L0
T = new_label() = L1
emit( T: )
t = generate(expr)=R2
emit( ifnot_goto t, E )
generate(statement)
emit( goto T )
emit( E: )

```

```

Code
L1:
R0 = load c
R1 = 20
R2 = R0 < R1
not_goto R2,L0
R3 = load c
R4 = 2
R5 = R3 + R2

```

```

r = generate(expr2)
l = lgenerate(expr2)
emit( store *l = r )
return r

```

```

t1 = generate(expr1)=R3
t2 = generate(expr2)=R4
r = new_temp() = R5
emit( r = t1 op t2 )
return r

```

Tufts University Computer Science

Example

```

E = new_label() = L0
T = new_label() = L1
emit( T: )
t = generate(expr)=R2
emit( ifnot_goto t, E )
generate(statement)
emit( goto T )
emit( E: )

```

```

Code
L1:
R0 = load c
R1 = 20
R2 = R0 < R1
not_goto R2,L0
R3 = load c
R4 = 2
R5 = R3 + R2

```

```

r = generate(expr2)=R5
l = lgenerate(expr2)
emit( store *l = r )
return r

```

Tufts University Computer Science

Example

```

E = new_label() = L0
T = new_label() = L1
emit( T: )
t = generate(expr)=R2
emit( ifnot_goto t, E )
generate(statement)
emit( goto T )
emit( E: )

```

```

Code
L1:
R0 = load c
R1 = 20
R2 = R0 < R1
not_goto R2,L0
R3 = load c
R4 = 2
R5 = R3 + R2
R6 = &c

```

```

r = generate(expr2)=R5
l = lgenerate(expr2)
emit( store *l = r )
return r

```

```

r = new_temp() = R6
emit( r = &v )
return r

```

Something like:
R6 = base + offset

Tufts University Computer Science

Example

```

E = new_label() = L0
T = new_label() = L1
emit( T: )
t = generate(expr)=R2
emit( ifnot_goto t, E )
generate(statement)
emit( goto T )
emit( E: )

```

```

Code
L1:
R0 = load c
R1 = 20
R2 = R0 < R1
not_goto R2,L0
R3 = load c
R4 = 2
R5 = R3 + R2
R6 = &c
store [R6]=R5

```

```

r = generate(expr2)=R5
l = lgenerate(expr2)=R6
emit( store *l = r )
return r

```

Tufts University Computer Science

Example

```

E = new_label() = L0
T = new_label() = L1
emit( T: )
t = generate(expr)=R2
emit( ifnot_goto t, E )
generate(statement)
emit( goto T )
emit( E: )
  
```

```

Code
L1:
R0 = load c
R1 = 20
R2 = R0 < R1
not_goto R2,L0
R3 = load c
R4 = 2
R5 = R3 + R2
R6 = & c
store [R6]=R5
goto L1
L0:
  
```

Tufts University Computer Science

Example

```

Code
R7 = load c
R8 = 0
R9 = R7 == R8
not_goto R9,L3
L1:
R0 = load c
R1 = 20
R2 = R0 < R1
not_goto R2,L0
R3 = load c
R4 = 2
R5 = R3 + R2
R6 = & c
store [R6]=R5
goto L1
L0:
goto L4
L3:
. . .
L4:
  
```

Tufts University Computer Science

Nesting

```

Code
R7 = load c
R8 = 0
R9 = R7 == R8
not_goto R9,L3
L1:
R0 = load c
R1 = 20
R2 = R0 < R1
not_goto R2,L0
R3 = load c
R4 = 2
R5 = R3 + R2
R6 = & c
store [R6]=R5
goto L1
L0:
goto L4
L3:
. . .
L4:
  
```

Tufts University Computer Science

Registers only

- Assume all variables are in registers
 - Add loads and stores later
 - Change *generate(v)*

```

r = new_temp()
emit( r = load v )
return r
  
```

```

Code
R7 = load c
R8 = 0
R9 = R7 == R8
not_goto R9,L3
L1:
R0 = load c
R1 = 20
R2 = R0 < R1
not_goto R2,L0
R3 = load c
R4 = 2
R5 = R3 + R2
R6 = & c
store [R6]=R5
goto L1
L0:
goto L4
L3:
. . .
L4:
  
```

Tufts University Computer Science

Registers only

- Proliferation of registers and labels

Many CPUs have a fast $c == 0$ test

Can use accumulators: $c = c + 2$

Label leads to another goto; may have multiple labels

```

Code
R7 = c
R8 = 0
R9 = R7 == R8
not_goto R9,L3
L1:
R0 = c
R1 = 20
R2 = R0 < R1
not_goto R2,L0
R3 = c
R4 = 2
R5 = R3 + R2
c = R5
goto L1
L0:
goto L4
L3:
. . .
L4:
  
```

Tufts University Computer Science

Efficient lowering

- Reduce number of temporary registers
 - Don't copy variable values unnecessarily
 - Accumulate values, when possible
 - Reuse temporaries, where possible
- Generate more efficient labels
 - Don't generate multiple adjacent labels
 - Avoid goto-label-goto

Tufts University Computer Science

Avoiding extra copies

- Basic algorithm
 - Recursive generation traverses to leaves
 - At leaves, generate: $R = v$ or $R = c$
- Improvement
 - Stop recursion one level early
 - Check to see if children are leaves
 - Don't call generate recursively on variables, constants



Avoiding copies

expr1 op expr2



```

if (expr1 is var)
    t1 = (Var) expr1
else
    t1 = generate(expr1)
if (expr2 is var)
    t2 = (Var) expr2
else
    t2 = generate(expr2)
r = new_temp()
emit( r = t1 op t2 )
return r
    
```



Example

- **Expr1** is (a+b)
 - Not a leaf
 - Recursively generate code
 - Return temp
- **Expr2** is c
 - Return c
- **Emit** ($R0 * c$)

(a + b) * c

Code

```

R0 = a + b
R1 = R0 * c
    
```



Use accumulation

- **Idea:**
 - We only need 2 registers to evaluate expr1 op expr2
 - Reuse temp assigned to one of the subexpressions

```

if (expr1 is var)
    t1 = (Var) expr1
else
    t1 = generate(expr1)
if (expr2 is var)
    t2 = (Var) expr2
else
    t2 = generate(expr2)
emit( t1 = t1 op t2 )
return t1
    
```



Example

- Combined:
 - Remove copies
 - Accumulate value
 - Only need one register
- (Original version might have required as many as 6)

(a + b) * c

Code

```

R0 = a + b
R0 = R0 * c
    
```



Reuse of temporaries

- **Idea:**
 - Can **generate**(expr1) and **generate**(expr2) share temporaries?
 - Yes, except for t1 and t2
 - Observation: temporaries have a limited lifetime
 - Lifetime confined to a subtree

expr1 op expr2



```

t1 = generate(expr1)
t2 = generate(expr2)
r = new_temp()
emit( r = t1 op t2 )
return r
    
```

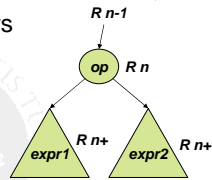


Reuse of temporaries

- Subtrees can share registers

- Algorithm:

- Use a stack of registers
- Start at $\# = 0$
- Each call to generate:
 - "Push" next number
 - Use any register $> \#$
 - When done, "pop" back up



```
R# = expr1  
R# = R# op expr2
```



Next time...

- Introduction to optimization
- Then: instruction selection

