Core ML

COMP 105 Assignment

Due Tuesday, October 16, 2018 at 11:59PM

Contents

Overview

Prelude

Setup

ML’s initial basis

Unit testing

Things you need to review before starting

What we expect from your submission

We expect the right types

We expect wise, well-formatted unit tests

We expect case analysis only when necessary.

We don’t expect written algebraic laws

We expect an acceptable style

We expect you to avoid forbidden functions and constructs

Reading comprehension (10%)

Programming problems to solve individually (75%)

Defining functions using clauses and patterns

Lists

Handling exceptions

Arithmetic by pattern matching on constructed data

An immutable, persistent alternative to linked lists

One problem you can do with a partner (15%)

Extra credit

Avoid common mistakes

What to submit and how to submit it

Submitting your individual work

Submitting your improved μScheme interpreter

How your work will be evaluated
Overview

This assignment gets you acclimated to programming with pattern matching and types. These features are found in such languages as Standard ML, Haskell, Elm, OCaml, Reason, F#, Scala, Idris, Agda, and Coq/Gallina. These languages excel at analysis and implementation of other languages; they are the technology of choice for applications like compilers, verification, and other static analysis, including of security properties. In the next few weeks, you will explore such applications by using Standard ML to implement type systems and lambda calculus.

The assignment has three parts:

- To begin, you will answer some questions about reading.
- On your own, you will write many small functions.
- Possibly working with a partner, you will make a small change to the μScheme interpreter that is written in ML (in Chapter 5).

After completing this assignment, you will be ready to tackle serious programming tasks in Standard ML.

Prelude

Setup

COMP 105 uses two different implementations of Standard ML. For the small problems, we recommend Moscow ML, which is in /usr/sup/bin/mosml. To start Moscow ML, use

```
ledit mosml -P full -I /comp/105/lib
```

If everything is working correctly, you should see this prompt:

```
Moscow ML version 2.10-3 (Tufts University, April 2012)
Enter ‘quit();’ to quit.
```

If you don’t see the Tufts name, send an immediate email of complaint to staff@cs.tufts.edu, with a copy to comp105-grades@cs.tufts.edu.

For the large problem, we recommend a native-code compiler called mlon (pronounced “Milton”).

ML’s initial basis

As in the hofs and continuations assignments, we expect you to use the initial basis, which is properly known as the Standard ML Basis Library. By the standards of popular languages, the basis is small, but it is still much more than you can learn in a week. Fortunately, you only have to learn a few key parts:

- Type constructors `list`, `option`, `bool`, `int`, `string`, and `order`

1 mailto:staff@cs.tufts.edu
2 http://www.sml-family.org/Basis/
3 http://sml-family.org/Basis/option.html#SIG:OPTION.option:TY
4 http://sml-family.org/Basis/general.html#SIG:GENERAL.order:TY
- Modules List\(^5\) and Option\(^6\), including List.filter, List.exists, List.find, and others
- Other module functions Int.toString, Int.compare, and String.compare
- Top-level functions o, print (for debugging), map, app, foldl, foldr
- Our own Unit module, which is not part of the Basis Library but is described in our ML Learning Guide\(^7\).

The most convenient guide to the basis is the Moscow ML help system; type

```
- help "";
```

at the mosml interactive prompt. The help file is badged incorrectly, but as far as I know, it is up to date.

If you have Jeff Ullman’s book, you need to know that Chapter 9 describes the 1997 basis, which is out
of date: today’s compilers use the 2004 basis, which is a standard. But there are only a few differences,
primarily in I/O and arrays. The most salient difference is in the interface to TextIO.inputLine\(^8\).

### Unit testing

Although Standard ML does not have check-expect and friends built in, we can do a lot with higher-
order functions. Here are some examples of tests written with our Unit module:

```ml
val () =
    Unit.checkExpectWith Int.toString "2 means the third"
    (fn () => List.nth ([1, 2, 3], 2))
    3

val () = (* this test fails *)
    Unit.checkExpectWith Bool.toString "2 is false"
    (fn () => List.nth ([true, false, true], 2))
    false

val () = Unit.reportWhenFailures ()
```

If you fire up mosml, load the Unit module, and paste in this code (followed by a semicolon), you’ll see
these reports:

```
$ mosml -P full -I /comp/105/lib
Moscow ML version 2.10-3 (Tufts University, April 2012)
Enter ‘quit();’ to quit.
- load "Unit";
> val it = () : unit
- val () = ... ;
In test '2 is false', expected value false but got true
One of two internal Unit tests passed.
```

\(^5\)http://sml-family.org/Basis/list.html#List:STR:SPEC
\(^6\)http://sml-family.org/Basis/option.html#Option:STR:SPEC
\(^7\)../readings/ml.html#unit-testing
\(^8\)http://sml-family.org/Basis/text-io.html#SIG:TEXT_IO.inputLine:VAL
You’ll use `Unit.checkExpectWith` to write your own unit tests. You’ll also use `Unit.checkAssert` and `Unit.checkExnWith`. The details are in the ML Learning Guide\(^9\).

**Things you need to review before starting**

We provide a guide to *Learning Standard ML*\(^10\). *Learning Standard ML* will guide you to other reading.

We also provide a handout on applying the 9-step design process to Standard ML\(^11\). This handout includes the key code excerpts needed to design and program with standard type constructors `list`, `option bool`, `int`, `string`, and `order`, as well as the `tree` constructor on the homework and the `exp` constructor from \(\mu\)Scheme.

**What we expect from your submission**

We expect you will submit code that compiles, has the types given in the assignment, is acceptably styled, is tested, avoids redundant case analysis, and avoids forbidden functions and constructs. Code that does not compile, that has the wrong types, or that uses forbidden functions or constructs will earn No Credit. Code that is untested or has redundant case analysis may earn disappointing grades for structure and organization.

**We expect the right types**

As always, your code is assessed in part through automated testing. To be testable, each function must not only have the correct *name*; it must also have the correct *type*. Your type definitions must also match the type definitions given in the assignment.

To help you check types, I provide a script that analyzes your solution and reports on any type mismatches that it finds. On the command line, run

```
mml-sanity-check warmup.sml
```

The “sanity check” reports both missing and ill-typed functions. Any function that does not pass the sanity check will earn No Credit. The sanity check looks at the types listed here:

\[
\begin{align*}
\text{val} & \text{ firstVowel : char list -> bool} \\
\text{val} & \text{ mynull : 'a list -> bool} \\
\text{val} & \text{ reverse : 'a list -> 'a list} \\
\text{val} & \text{ minlist : int list -> int} \\
\text{exception} & \text{Mismatch} \\
\text{val} & \text{ zip : 'a list * 'b list -> (a * b) list} \\
\text{val} & \text{ ziptoo : 'a list * 'b list -> (a * b) list} \\
\text{val} & \text{ pairfoldrEq : ('a * 'b * 'c -> 'c) -> c -> 'a list * 'b list -> c} \\
\text{val} & \text{ concat : 'a list list -> 'a list}
\end{align*}
\]

\(^9\)../readings/ml.html#unit-testing
\(^10\)../readings/ml.pdf
\(^11\)../handouts/mlproofs.pdf
We expect wise, well-formatted unit tests

By this time, we expect that you understand the value of unit tests. Grading will focus on your code; except where specifically requested below (natural-number arithmetic, free-variable analysis), your unit tests won’t be graded. But we still expect the following:

- You will **indent all unit tests by eight spaces**. This indentation will enable graders to focus on your code.
- You will use unit tests wisely. If a function is simple, do take a minute to validate it with a unit test. If a function is not so simple, develop unit tests in the same way you have done for the past three assignments: one unit test per case in the code.
• If you need debugging help during office hours, we expect that your code will be accompanied by failing unit tests. (If you cannot get your code to typecheck, we will help you do this without unit tests. But if you need help getting code to produce right answers, we will demand to see your unit tests.)

We expect case analysis only when necessary.

Case analysis is the enemy. All the more so when it is not necessary. Redundant case analysis is a problem in all levels of programming, but as you are learning ML, it is especially easy to fall into. Redundant case analysis typically manifests in one of two ways:

1) Two cases are present in a `fun`, or case, but one is completely subsumed by the other. The most common example is one case to handle the empty list and another case that handles all lists. The empty-list case is often redundant.

Example:

```haskell
fun append ([], ys) = ys
| append (xs, ys) = foldr (\) :: ys xs
```

In this code, the first case is subsumed by the second. It can be eliminated without changing the meaning of the code, and eliminating it typically improves performance.

2) A case analysis is performed where no case analysis is needed.

```haskell
fun sum [] = 0
| sum (n :: ns) = foldl (\) + n ns
```

These two cases should be replaced by a single case:

```haskell
fun sum ns = foldl (\) + 0 ns
```

We expect you to examine your code carefully and to remove all redundant case analyses.

We don’t expect written algebraic laws

We expect you to continue using a systematic design process, but because ML code is so close to algebraic laws, we don’t expect you to write algebraic laws separately. If you come to office hours, however, we do expect you to be able to talk about algebraic laws and to write them on the board.

We expect an acceptable style

Nobody can learn good style in a week. But you can learn to imitate somebody else’s style, and we expect you to be judicious about what style you imitate. You have access to books by Ullman, Ramsey, and Harper, and to a technical report by Toffe. These sources are not equally good:

• Ullman provides the most gentle introduction to ML, and he provides the most information about ML. His book is especially good for programmers whose primary experience is in C-like languages. But, to put it politely, Ullman’s code is not idiomatic. Much of what you see from Ullman should not be imitated.
• Ramsey’s code, starting in Chapter 5, is a better guide to what ML should look like. Harper’s code is also very good, and Tofte’s code is reasonable.

On this assignment, we expect you to devote a little effort to good style. Focus on getting your code working first. Then submit it. Then pick up our “Style Guide for Standard ML Programmers”\textsuperscript{12}, which contains many examples of good and bad style. Edit your code lightly to conform to the style guide, and submit it again.

In the long run, we expect you to master and follow the guidelines in the style guide\textsuperscript{13}.

**We expect you to avoid forbidden functions and constructs**

While not everybody can learn good style quickly, everybody can learn to avoid the worst faults. In ML, you must avoid these functions and idioms:

• Unlike μScheme, Standard ML provides a \texttt{length} function in the initial basis. It is banned. The entire assignment must be solved without using \texttt{length}.

**Solutions that use \texttt{Length} will earn No Credit.**

• Use function definition by pattern matching. Do not use the functions \texttt{null}, \texttt{mynull}, \texttt{hd}, and \texttt{tl}; use patterns instead.

**Solutions that use \texttt{hd} or \texttt{tl} will earn No Credit.**

• Except for functions given below, do not define auxiliary functions at top level. Use \texttt{local} or \texttt{let}. You will find it useful to use \texttt{local} to define functions for use in unit tests.

**Solutions that define auxiliary functions at top level will earn No Credit.**

• Do not use \texttt{open}; if needed, use short abbreviations for common structures. For example, if you want frequent access to the \texttt{ListPair} structure, you can write

```ml
structure LP = ListPair
```

and from there on you can refer to, e.g., \texttt{LP.map}.

**Solutions that use \texttt{open} may earn No Credit for your entire assignment.**

• Unless the problem explicitly says it is OK, do not use any imperative features.

Unless explicitly exempted, **solutions that use imperative features will earn No Credit.**

**Reading comprehension (10%)**

These problems will help guide you through the reading. We recommend that you complete them before starting the other problems below. You can download the questions\textsuperscript{14}.

\textsuperscript{12}../handouts/mlstyle.pdf
\textsuperscript{13}../handouts/mlstyle.pdf
\textsuperscript{14}/cqs.ml.txt
1. Read section 5.1 of Harper\textsuperscript{15} about tuple types and tuple patterns. Also look at the list examples in sections 9.1 and 9.2 of Harper.

Now consider the pattern \((x::y::zs, w)\). For each of the following expressions, tell whether the pattern matches the value denoted. If the pattern matches, say what values are bound to the four variables \(x, y, zs,\) and \(w\). If it does not match, explain why not.

(a) \(([1, 2, 3], ("COMP", 105))\)
(b) \(("COMP", 105), [1, 2, 3])\)
(c) \(("COMP", 105), (1, 2, 3))\)
(d) \("COMP", "105", true)\)
(e) \((\{true, false\}, 2.718281828)\)

Answers here:
(a) 
(b) 
(c) 
(d) 
(e) 

You are starting to be ready to use pattern matching.

2. Read the descriptions of patterns and example values (steps 3 and 2) in the handout “Program Design with ML Types and Pattern Matching\textsuperscript{16}.” In the same handout, look at Table 1, including the Types of parts column. Using the ideas you find there, prepare to answer questions about this expression:

\[
\text{case } f \ (x, y, z) \ \\
\quad \text{of } [] \quad \Rightarrow \text{raise Empty} \\
\quad | \ w :: ws \Rightarrow \text{if } p \ w \text{ then SOME } w \text{ else NONE}
\]

You are told that the subexpression \(f \ (x, y, z)\) has type \(\text{`a list}\). Using that information, give the type of each of these code fragments, which are built from parts of patterns:

(a) The type of the pattern \(w :: ws\)

\[
\rightarrow
\]

(b) The type of the variable \(ws\)

\[
\rightarrow
\]

(c) The type of the expression \(\text{SOME } w\)

\[
\rightarrow
\]

You are ready to use pattern matching on common ML types.

\textsuperscript{15}http://www.cs.cmu.edu/~rwh/isml/book.pdf
\textsuperscript{16}../handouts/mlproofs.pdf
3. Read the section on unit testing in the guide to learning ML\textsuperscript{17}. Read about infix function names in step 3 of the “design steps” section of the handout “Program Design with ML Types and Pattern Matching”\textsuperscript{18}. And read the section on unit testing in this homework.

Now, using the interpreter to be sure your answer is well typed, translate the following failing unit test into ML:

\[
(\text{check-expect } (\text{foldl } + \ 0 \ '(1 \ 2 \ 3)) \ 7)
\]

Your translation:

→

You are ready to write unit tests.

4. In chapter 5 of \textit{Programming Languages: Build, Prove, and Compare}, the \texttt{eval} code for applying a function appears in code chunk 371d. In evaluating \texttt{APPLY (f, args)}, if expression \texttt{f} does not evaluate to either a primitive function or a closure, the code raises the \texttt{RuntimeException} exception.

(a) Show a piece of \muScheme code that would, when evaluated, cause chunk 371d to raise the \texttt{RuntimeException} exception. \textbf{We are asking for \muScheme code that the interpreter runs, not for ML code.}

→

(b) When exception \texttt{RuntimeException} is raised, what happens from the user’s point of view? That is, if a person is interacting with the \muScheme interpreter at the read/eval/print loop, and evaluating an expression raises the \texttt{RuntimeException} exception, what user-visible actions will the interpreter take?

You are ready to write \texttt{zip} and to write environment functions that use exceptions.

5. “Free” variables are those that are not bound to a value in the current scope. You can find a longer discussion and precise definition in section 5.11 of \textit{Build, Prove, and Compare}, which starts on page 382. Read the section and identify the free variables of the expressions below. \textit{Remember that a name occurring in function position can be a free variable.}

(a) Free variables of \((\text{lambda} (x) \ (\text{lambda} (y) (\text{equal?} \ x \ y)))\)

→

(b) Free variables of \((\text{lambda} (y) (\text{equal?} \ x \ y)))\)

→

(c) Free variables of

\[
(\text{lambda} (s1 \ s2)
  (\text{if} \ (\text{or} \ (\text{atom?} \ s1) \ (\text{atom?} \ s2))
    (= s1 \ s2)
    (\text{and} \ (\text{equal?} \ (\text{car} \ s1) \ (\text{car} \ s2))
      (\text{equal?} \ (\text{cdr} \ s1) \ (\text{cdr} \ s2)))))
\]

→

\textsuperscript{17}../readings/ml.pdf
\textsuperscript{18}../handouts/mlproofs.pdf
You are ready to improve the μScheme interpreter. You may do this with a partner. You and your partner will turn your answers to parts (a) and (b) into unit tests.

Programming problems to solve individually (75%)

Working on your own, please solve the problems below. Place your solutions in file warmup.sml. At the start of each problem, please place a short comment, like

(***** Problem A *****)

At the very end of your warmup.sml, please place the following line:

val () = Unit.reportWhenFailures () (* put me at the _end_ *)

This placement will ensure that if a unit test fails, you are alerted.

To receive credit, your warmup.sml file must compile and execute in the Moscow ML system, and your functions must have the right types. At minimum, your code must compile without warnings or errors, and it must pass our sanity tests:

% /usr/sup/bin/mosmlc -toplevel -I /comp/105/lib -c warmup.sml
% ml-sanity-check warmup.sml

Defining functions using clauses and patterns

Related Reading for problems A and B: In Learning Standard ML\(^{19}\) read about Expressions (sections I, II, and III), Data (I, II, and II), Inexhaustive pattern matches, Types (I), Definitions (III, IV), and Expressions (VIII).

A. Define a function mynull : 'a list -> bool, which when applied to a list tells whether the list is empty. Avoid if, and make sure the function takes constant time. Do not use any functions from the Standard Basis. Make sure your function has the same type as the null in the Standard Basis.

B. Define a function firstVowel : char list -> bool that takes a list of lower-case letters and returns true if the first character is a vowel (aeiou) and false if the first character is not a vowel or if the list is empty. Use the wildcard symbol _ whenever possible, and avoid if.

Lists

Related Reading for problems C to F: In Learning Standard ML\(^{20}\), in addition to the section noted above, read about Types (III), and Exceptions. You will need to understand lists and pattern matching on lists (see Data III). You may also wish to read the section on Curried Functions.

C. Functions foldl and foldr are predefined with type

('a * 'b -> 'b) -> 'b -> 'a list -> 'b

They are like the μScheme versions except the ML versions are Curried.

---

\(^{19}\) ../readings/ml.html
\(^{20}\) ../readings/ml.html
1. Define reverse : ‘a list -> ‘a list using foldl or foldr.
   In ML, the reverse function is in the initial basis as rev.

2. Implement minlist : int list -> int, which returns the smallest element of a nonempty list of integers. Use foldl or foldr.
   If given an empty list of integers, your solution must fail (e.g., by raise Match).
   Your solution should work regardless of the representation of integers: it should not matter how many bits are used to represent a value of type int. (Hint: The course solution max* from the hofs\textsuperscript{21} homework works regardless of the representation of integers. Perhaps you can steal an idea from it.)

Do not use recursion in either part of this problem.

D. Define a function zip : ‘a list * ‘b list -> (‘a * ‘b) list that takes a pair of lists (of equal length) and returns the equivalent list of pairs. If the lengths don’t match, raise the exception Mismatch, which you must define. Do not use any functions from the Standard Basis Library.

You are welcome to translate a solution from μScheme, but you must either use a clausal definition or write code containing at most one case expression. Do not use if.

E. Define a function
   \[
   \text{val pairfoldrEq : (‘a * ‘b * ‘c -> ‘c) -> ‘c -> ‘a list * ‘b list -> ‘c}
   \]
   that applies a three-argument function to a pair of lists of equal length, using the same order as foldr.
   Do not use any functions from the Standard Basis Library. If pairfoldrEq’s contract is violated (by calling it on lists of unequal lengths), it must raise an exception.

Define a function ziptoo : ‘a list * ‘b list -> (‘a * ‘b) list which does exactly the same things as zip but which uses pairfoldrEq for its implementation.

F. Define a function
   \[
   \text{val concat : ‘a list list -> ‘a list}
   \]
   which takes a list of lists of ‘a and produces a single list of ‘a containing all the elements in the correct order. For example,

   - concat [[1], [2, 3, 4], [], [5, 6]];  
   > val it = [1, 2, 3, 4, 5, 6] : int list

Do not use if. You may use functions from the Standard Basis Library, except for List.concat—code that uses List.concat will earn No Credit.

To get full credit for this problem, your function should use no unnecessary cons cells. Keep in mind the cost of appending two lists.

**Handling exceptions**

**Related Reading** for problem G: In *Learning Standard ML*\textsuperscript{22}, read the section on Curried functions. Read the sections on Types (III) and Data (IV). Make sure you understand the difference between type

\textsuperscript{21}../hofs.html
\textsuperscript{22}../readings/ml.html
abbreviations (type) and algebraic data types (datatype)—both are called “types.” Read the section on Exceptions, and make sure you know both how to raise and how to handle an exception.

G. Environments with exceptions.
(For this problem, it is OK to use imperative features.)

Define type ‘a env as an abbreviation for a function type, and define exception NotFound, both as follows:

type 'a env = string -> 'a
exception NotFound of string

Now define these functions:

val emptyEnv : 'a env = (* ... *)
val bindVar : string * 'a * 'a env -> 'a env = (* ... *)
val lookup : string * 'a env -> 'a = (* ... *)

such that you can use ‘a env for a type environment or a value environment. On an attempt to look up an identifier that doesn’t exist, raise the exception NotFound. Don’t worry about efficiency.

Now define a function

val isBound : string * 'a env -> bool

that tells if a name is bound in the environment. Calling isBound must never raise an exception.

Hint: Review the “sets as characteristic functions” from the hofs assignment.

Arithmetic by pattern matching on constructed data

Languages like C and C++ enable you to do arithmetic only on as many bits as are in a machine word. More civilized languages allow arithmetic on as many bits as will fit in memory. Every computer scientist should know how this feature is implemented. In this assignment, we’ll implement arithmetic on natural numbers only, where a natural number is represented as a list of digits.

Our representation is based on the Decimal proof system from the proof-systems handout. To make the rules readable on the web and not just in the printout, I recapitulate them in informal English:

- A digit is an integer in the range 0 to 9 inclusive.
- Zero is a natural number (rule DecimalZero).
- If \( m \) is a natural number and \( d \) is a digit, then \( 10 \times m + d \) is a natural number (rule DecimalNat).

I choose to represent a natural number as a data structure defined by an algebraic data type:

```
datatype nat = ZERO
  | TIMES10PLUS of nat * int
```

(This data structure is equivalent to a list of digits with the least-significant digit first.) The meaning of the data structure is determined by an abstraction function and representation invariants:

- The abstraction function says that ZERO stands for the natural number zero, and if \( m \) is a natural number and \( d \) is a digit, then \( \text{TIMES10PLUS} \ (m, d) \) stands for the natural number \( m \times 10 + d \).

---

23 By this criterion, Standard ML is not civilized, because the so-called “large-integer arithmetic” is optional.
24 ../handouts/natproofs.pdf
• The representation invariants say that in any value of the form TIMES10PLUS \( m, d \),
  
  - \( m \) and \( d \) are not both zero.
  
  - \( d \) is a machine integer in the range \( 0 \leq d < 10 \).

The first representation invariant is maintained by the following function, which is called a “smart constructor”:

\[
\text{fun times10plus (ZERO, 0) = ZERO} \\
| \text{times10plus (m, d) = TIMES10PLUS (m, d)}
\]

If you like, you can add code to the second case to enforce the invariant on \( d \).

Here are a couple of useful special cases:

\[
(* \text{times10 : nat -> nat *}) \\
\text{fun times10 n = times10plus (n, 0)}
\]

\[
(* \text{natOfDigit : int -> nat *}) \\
\text{fun natOfDigit d = times10plus (ZERO, d)}
\]

\text{Put the definitions of nat, times10plus, and the special-case functions into your code.}

You will define functions for conversion, addition, subtraction. Multiplication is extra credit.

For testing, you will find it useful to convert a list of decimal digits to a natural number. You are welcome to use this code:

\[
\text{fun flip f (x, y) = f (y, x)} \\
(* \text{natOfDigits : int list -> nat *}) \\
\text{fun natOfDigits ds = foldl (flip times10plus) ZERO ds}
\]

\text{Related Reading:}

- If needed, review the reading on pattern matching you’ve already done.

- In Learning Standard ML\(^\text{25}\), read the section on datatypes—Data IV. Make sure you understand how to pattern match on constructed values.

- For detailed descriptions of algorithms for addition and subtraction, with examples, read Programming Languages: Build, Prove, and Compare, section 9.10.2, which starts on page 729. Chapter 9 is not fully included in your abridged edition, but the pages containing this section are there.

- To understand how op is used in the unit-test examples, consult Expressions VII: Infix operators as functions\(^\text{26}\) in Learning Standard ML.

\text{H. Natural-number conversions.}

You will convert between natural numbers, machine integers, and strings.

1. Define a function

\[
\text{val int0fNat : nat -> int}
\]

\(^{25}\)../readings/ml.html

\(^{26}\)../readings/ml.html#expressions-vii-infix-operators-as-functions
that converts a natural number into a machine integer, or if the natural number is too large, raises \texttt{Overflow}. (Use the built-in operators \texttt{+} and \texttt{*}, which do machine arithmetic and which automatically raise \texttt{Overflow} when needed.)

Example:

```markdown
- intOfNat (natOfDigits [1, 2, 3]);
> val it = 123 : int
```

2. Write a unit test confirming what the example shows: that \texttt{intOfNat (natOfDigits [1, 2, 3])} is 123.

3. Define a function

```
val natOfInt : int -> nat
```

that converts a nonnegative machine integer into a natural number.

Example:

```markdown
- natOfInt 2018;
> val it = TIMES10PLUS(TIMES10PLUS(TIMES10PLUS(TIMES10PLUS(ZERO, 2), 0), 1), 8) : nat
```

Use pattern matching, not if.

A nonnegative machine integer is either zero or it has the form \( n = 10 \times m + d \). In the second case, \( d \) is \((n \mod 10)\) and \( m \) is \((n \div 10)\).

4. Write a unit test confirming the \texttt{natOfInt} example.

5. Define function \texttt{natString}, which converts a \texttt{nat} to a string the way we normally write it (with the most significant digit first).

```
val natString : nat -> string
```

Examples:

```markdown
- natString (natOfDigits [3, 2, 1]);
> val it = "321" : string
- natString (natOfDigits [2, 0, 1, 8]);
> val it = "2018" : string
```

Function \texttt{natString} must never return an empty string.

Hint: Go back to the handout on proof systems for natural numbers. The representation above is based on the Decimal system. But the \texttt{natString} function needs to be based on the DecNumeral system. Find a way to deal with the difference.

6. Write a unit test confirming the \texttt{natString} example.

My solutions take 7 lines of code and 18 lines of (paranoid) unit tests.

I. \textit{Natural-number arithmetic}.

You will add and subtract natural numbers.

1. Define function \texttt{carryIntoNat : nat * int -> nat}. This function takes a natural number \( n \) and a carry bit \( c \), and it returns \( n + c \). A carry bit is a machine integer that is either 0 or 1.
The function is defined by these algebraic laws:

\[
\begin{align*}
carryIntoNat(n, 0) &= n \\
carryIntoNat(0, c) &= c \\
carryIntoNat(10 \cdot m + d, 1) &= \\
& \quad 10 \cdot carryIntoNat(m, (d + 1) \div 10) \\
& \quad + ((d + 1) \mod 10)
\end{align*}
\]

To convert these laws into code, you will need to write the natural-number patterns for 0 and for \(10 \cdot m + d\) as constructed-data patterns \textsc{Zero} and \textsc{Times10Plus} \((m, d)\). And you will need to write the natural-number arithmetic on the right-hand side using the smart constructor \textsc{times10plus}.

2. Define function \texttt{addWithCarry : nat * nat * int -> nat}. This function takes two natural numbers \(n_1\) and \(n_2\), and a carry bit \(c\), and it returns \(n_1 + n_2 + c\). To earn a passing grade, it must be capable of adding 30-digit numbers, regardless of the number of bits available in a machine integer.

The function is defined by these algebraic laws:

\[
\begin{align*}
\text{addWithCarry}(n_1, 0, c) &= \text{carryIntoNat}(n_1, c) \\
\text{addWithCarry}(0, n_2, c) &= \text{carryIntoNat}(n_2, c) \\
\text{addWithCarry}(10 \cdot m_1 + d_1, 10 \cdot m_2 + d_2, c) &= \\
& \quad \text{let } d = (d_1 + d_2 + c) \mod 10 \\
& \quad \text{val } c' = (d_1 + d_2 + c) \div 10 (* \text{ the "carry out" } *) \\
& \quad \text{in } 10 \cdot \text{addWithCarry}(m_1, m_2, c') + d \\
\end{align*}
\]

To convert these laws into code, you will need to write the natural-number patterns as constructed-data patterns, and you will need to write the final operation between \texttt{in} ... \texttt{end} using the smart-constructor function.

3. Define function \texttt{addNats : nat * nat -> nat}, as follows:

\[
\text{fun addNats}(n_1, n_2) = \text{addWithCarry}(n_1, n_2, 0)
\]

4. Define function \texttt{borrowFromNat : nat * int -> nat}. This function takes a natural number \(n\) and a borrow bit \(b\), and it returns \(n - b\), provided that \(n - b\) is a natural number. If \(n - b\) is not a natural number, \texttt{borrowFromNat} raises the exception \texttt{Negative}, which you will need to define. The borrow bit, like a carry bit, is a machine integer that is either 0 or 1.

The function is defined by these algebraic laws:

\[
\begin{align*}
\text{borrowFromNat}(n, 0) &= n \\
\text{borrowFromNat}(10 \cdot m + 0, 1) &= 10 \cdot \text{borrowFromNat}(m, 1) + 9 \\
\text{borrowFromNat}(10 \cdot m + d, 1) &= 10 \cdot m + (d - 1), \text{ where } d > 0
\end{align*}
\]

Notice there is no law for the left-hand side \texttt{borrowFromNat} \((0, 1)\). That’s because \(0 - 1\) is not a natural number—so if your code encounters this case, it should raise the \texttt{Negative} exception.

To convert these laws into code, you will need to write the natural-number patterns as constructed-data patterns, and you will need to write some of the arithmetic on the right-hand side using the smart-constructor function.

5. Define function \texttt{subWithBorrow : nat * nat * int -> nat}. This function takes two natural numbers \(n_1\) and \(n_2\), and a borrow bit \(b\), and if \(n_1 - n_2 - b\) is a natural number, it returns \(n_1 - n_2 - b\).
Otherwise it raises the Negative exception.

Like addWithCarry, subWithBorrow must be capable of subtracting 30-digit numbers.

The function is defined by these algebraic laws:

\[
\text{subWithBorrow} \ (n_1, 0, b) = \text{borrowFromNat} \ (n_1, b) \\
\text{subWithBorrow} \ (10 \times m_1 + d_1, 10 \times m_2 + d_2, b) = \\
\quad \text{let val } d = (d_1 - d_2 - b) \mod 10 \\
\quad \quad \text{val } b' = \text{if } d_1 - d_2 - b < 0 \text{ then } 1 \text{ else } 0 \text{ (* the "borrow out" *)} \\
\quad \quad \text{in } 10 \times \text{subWithBorrow} \ (m_1, m_2, b') + d \\
\text{end}
\]

**Alert:** These laws assume the Standard ML definition of \( \text{mod} \), which is not what you get from the hardware. The result of \( k \mod 10 \) is always nonnegative.

To convert these laws into code, you will need to write the natural-number patterns as constructed-data patterns, and you will need to write the final operation between \( \text{in } \ldots \text{ end} \) using the smart-constructor function.

6. Define function \( \text{subNats} : \text{nat} * \text{nat} \rightarrow \text{nat} \), as follows:

\[
\text{fun subNats} \ (n_1, n_2) = \text{subWithBorrow} \ (n_1, n_2, 0)
\]

Here is a unit test to confirm that subtracting too large a number raises the proper exception: it should raise Negative and not Match:

\[
\text{val () = } \\
\quad \text{Unit.checkExnSatisfiesWith natString "1 - 5"} \\
\quad \quad \text{(fn () => subNats (natOfDigits [1], natOfDigits [5]))} \\
\quad \quad \text("Negative", fn Negative => true | _ => false)
\]

If you trust your conversion functions from the previous problem, you can write unit tests using higher-order functions. Here is an example:

\[
\text{fun opsAgree name intop natop n1 n2 = } \\
\quad \text{Unit.checkExpectWith Int.toString name} \\
\quad \quad \text{(fn () => intOfNat (natop (natOfInt n1, natOfInt n2)))} \\
\quad \quad \text{(intop (n1, n2) handle Overflow => 0)}
\]

This function has type

\[
\text{val opsAgree} : \\
\quad \text{string -> (int * int -> int) -> (nat * nat -> nat) ->} \\
\quad \quad \text{int -> int -> unit}
\]

And it is used as follows

\[
\text{val () = opsAgree "123 + 2018" (op +) addNats 123 2018} \\
\text{val () = opsAgree "2018 - 123" (op -) subNats 2018 123} \\
\text{val () = opsAgree "2018 * 123" (op *) mulNats 2018 123} \\
\text{val () = opsAgree "100 - 1 " (op -) subNats 100 1}
\]

(Multiplication is for extra credit.)
My addition functions total 14 lines of code, not counting unit tests. My subtraction functions also total 14 lines of code, not counting unit tests.

**Hints:**

- Exploit the representation invariant. If a natural number matches the pattern TIMES10PLUS (\_, \_), the representation invariant guarantees that the number is not zero.
- To maintain the representation invariant, use TIMES10PLUS only in pattern matching. On the right-hand side of any algebraic law, build natural numbers using the smart constructor times10plus.

**An immutable, persistent alternative to linked lists**

**Related Reading** for problem J: In *Learning Standard ML*\(^{27}\), read the section on datatypes—Data IV. Make sure you understand how to pattern match on constructed values.

J. For this problem I am asking you to define your own representation of a new abstraction: the list with finger. A list with finger is a nonempty sequence of values, together with a “finger” that points at one position in the sequence. The abstraction provides constant-time insertion and deletion at the finger.

**This is a challenge problem.** The other problems on the homework all involve old wine in new bottles. To solve this problem, you have to *think* of something new.

1. Define a representation for type `'a flist`. (Before you can define a representation, you will want to study the rest of the parts of this problem, plus the test cases.)

   **Document** your representation by saying, in a short comment, what sequence is meant by any value of type `'a flist`.

2. Define function

   ```ml
   val singletonOf : 'a -> 'a flist
   ```

   which returns a sequence containing a single value, whose finger points at that value.

3. Define function

   ```ml
   val atFinger : 'a flist -> 'a
   ```

   which returns the value that the finger points at.

4. Define functions

   ```ml
   val fingerLeft : 'a flist -> 'a flist
   val fingerRight : 'a flist -> 'a flist
   ```

   Calling `fingerLeft xs` creates a new list that is like `xs`, except the finger is moved one position to the left. If the finger belonging to `xs` already points to the leftmost position, then `fingerLeft xs` should raise the same exception that the Basis Library raises for array access out of bounds. Function `fingerRight` is similar. Both functions must run in **constant time and space**.

   Please think of these functions as “moving the finger”, but remember **no mutation is involved**.

   Instead of changing an existing list, each function creates a new list.

5. Define functions

\(^{27}\)../readings/ml.html
Val `deleteLeft : 'a flist -> 'a flist
Val `deleteRight : 'a flist -> 'a flist

Calling `deleteLeft xs` creates a new list that is like xs, except the value x to the left of the finger
has been removed. If the finger points to the leftmost position, then `deleteLeft` should raise the
same exception that the Basis Library raises for array access out of bounds. Function `deleteRight`
is similar. Both functions must run in **constant time and space**. As before, no mutation is involved.

6. Define functions

Val `insertLeft : 'a * 'a flist -> 'a flist
Val `insertRight : 'a * 'a flist -> 'a flist

Calling `insertLeft (x, xs)` creates a new list that is like xs, except the value x is inserted
to the left of the finger. Function `insertRight` is similar. Both functions must run in **constant
time and space**. As before, no mutation is involved. (These functions are related to “cons”.)

7. Define functions

Val `ffoldl : ('a * 'b -> 'b) -> 'b -> 'a flist -> 'b
Val `ffoldr : ('a * 'b -> 'b) -> 'b -> 'a flist -> 'b

which do the same thing as `foldl` and `foldr`, but ignore the position of the finger.

Here is a simple test case, which should produce a list containing the numbers 1 through 5 in order. You
can use `ffoldr` to confirm.

Val `test = singletonOf 3
Val `test = insertLeft (1, test)
Val `test = insertLeft (2, test)
Val `test = insertRight (4, test)
Val `test = fingerRight test
Val `test = insertRight (5, test)

You’ll want to test the delete functions as well.

Here are some properties you can use for testing:

Val `deleteLeft (insertLeft (x, xs)) == xs
Val `fingerLeft (insertLeft (x, xs)) ==
    fingerRight (insertRight (x, fingerLeft xs))

Properties like these can be converted to functions:

Val `prop1 = fn (x, xs) => deleteLeft (insertLeft (x, xs)) = xs
Val `prop2 = fn (x, xs) => fingerLeft (insertLeft (x, xs)) =
    fingerRight (insertRight (x, fingerLeft xs))

Property `prop1` holds of any list with finger. Property `prop2` holds only if it is possible to move the finger
to the left—otherwise it raises an exception.

Hints: The key is to come up with a good representation for “list with finger.” Once you have a good
representation, the code is easy: over half the functions can be implemented in one line each, and no func-
tion requires more than two lines of code. The biggest risk is that you copy/paste something incorrectly
for left/right, and you can mitigate this risk through testing.
One problem you can do with a partner (15%)

As noted above, ML-like languages excel at analyzing and manipulating programs. In the coming month, you will write many functions that analyze expressions; this problem will get you off to a good start.

The problem is numbered 2 because that’s the problem number in the book. You won’t be doing exercise 1, so you’re not missing anything.

**Related Reading** for exercise 2: *Build, Prove, and Compare*, section 5.11, which starts on page 382. Focus on the proof system for judgment $y \in \mathit{fv}(e)$; it is provable exactly when $\mathit{freeIn} \ e \ y$, where $\mathit{freeIn}$ is the most important function in exercise 2. Also read function $\mathit{eval}$ in section 5.4. You will modify the case for evaluating $\text{LAMBDA}$.

2. **Improving closures.**

When a compiler translates a lambda expression, it doesn’t store the entire environment in the closure; it stores only the free variables of the lambda expression. Interpreters for languages like Lua and Python work in the same way. You’ll implement this code improvement in an interpreter for µScheme.

This problem appears in *Build, Prove, and Compare* as exercise 2 on page 389, and you’ll solve it in a prelude and four parts:

- The prelude is to go to your copy of the book code and copy the file bare/uscheme-ml/mlscheme.sml to your working directory. (This code contains all of the interpreter from Chapter 5.) Then make another copy and name it mlscheme-improved.sml. You will edit mlscheme-improved.sml.
- The first part is to implement the free-variable predicate
  
  ```
  val freeIn : exp -> name -> bool.
  ```

  This predicate tells when a variable appears free in an expression. It implements the proof rules in section 5.11 of the book, which starts on page 382.

  During this part I recommend that you **compile early and often** using

  ```
  /usr/sup/bin/mosmlc -c -toplevel -l /comp/105/lib/mlscheme-improved.sml
  ```

  We also require **unit tests** for $\mathit{freeIn}$. At minimum, write two tests for each short example in the reading-comprehension questions: one for a variable that is free, and one for a variable that appears in the expression but is not free. Unit tests for LET forms are recommended but not required.

- The second part is to write a function that takes a pair consisting of a $\text{LAMBDA}$ body and an environment, and returns a better pair containing the same $\text{LAMBDA}$ body paired with an environment that contains only the free variables of the $\text{LAMBDA}$. (In the book, in exercise 1 starting on page 388, this environment is explained as the *restriction* of the environment to the free variables.) I recommend that you call this function $\mathit{improve}$, and that you give it the type

  ```
  val improve : (name list * exp) * 'a env -> (name list * exp) * 'a env
  ```

- The third part is to use $\mathit{improve}$ in the evaluation case for $\text{LAMBDA}$, which appears in the book on page 371c. You simply apply $\mathit{improve}$ to the pair that is already there, so your improved interpreter looks like this:

  ```
  (* more alternatives for \[ev\] for \uscheme 371c *)
  | ev (LAMBDA (xs, e)) = ( errorIfDups ("formal parameter", xs, "lambda")
  ; CLOSURE (improve ((xs, e), rho))
  ```
The fourth and final part is to see if it makes a difference. You will compile both versions of the μScheme interpreter using MLton, which is an optimizing, native-code compiler for Standard ML. The compiler requires some annoying bureaucracy, but it compensates by providing native-code speeds.

The original file, which has no unit tests, can be compiled without bureaucracy:

```
mlton -verbose 1 -output mlscheme mlscheme.sml
```

(If plain `mlton` doesn’t work, try `/usr/sup/bin/mlton`.)

Compiling your improved version requires some bureaucracy to incorporate the `Unit` module.

```
copy-105-ml-files-here
mlton -verbose 1 -output mlscheme-improved mlscheme-with-unit.mlb
```

The file `mlscheme-with-unit.mlb` tells MLton to compile your code with our `Unit` module. If you wish to do this on your own computer, you will also need files `unit.mlb` and `unit-mlton.sml` from `/comp/105/lib`, and you will have to edit `mlscheme-with-unit.mlb` to refer to your local copy of `unit.mlb`, not the one in `/comp/105/lib`.

Once compiled, you will run both versions and see if the “improvement” is measurable. For measurement, I have provided a script you can use. I also recommend that you compare the performance of the ML code with the performance of the C code in the course directory.

To get a good measurement, you will need to turn off the “CPU throttling” feature that is built into our interpreters. Use the following arcane Unix commands:

```
env BPCOPTIONS=nothrottle time run-exponential-arg-max 22 ./mlscheme
env BPCOPTIONS=nothrottle time run-exponential-arg-max 22 ./mlscheme-improved
env BPCOPTIONS=nothrottle time run-exponential-arg-max 22 /comp/105/bin/uscheme
```

(If you get an error message along the lines of “CPU time exhausted,” something is wrong.)

**Hints:**

- Focus on function `freeIn`. This is the only recursive function and the only function that requires case analysis on expressions. And it is the only function that requires you to understand the concept of free variables. All of these concepts are needed for future assignments.

  Understanding free variables is hard, but once you understand, the coding is easy.

- In Standard ML, the μScheme function `exists?` is called `List.exists`. You’ll have lots of opportunities to use it. If you don’t use it, you’re making extra work for yourself.

In addition to `List.exists`, you may find uses for `map`, `foldr`, `foldl`, or `List.filter`.

You might also find a use for these functions, which are already defined for you:

```ml
fun fst (x, y) = x
fun snd (x, y) = y

fun member (y, []) = false
| member (y, z::zs) = y = z orelse member (y, zs)
```
• The case for LETSTAR is gnarly, and writing it adds little to the experience. Here are two algebraic laws which may help:

\[
\text{freeIn (LETX (LETSTAR, [], e)) y = freeIn e y}
\]

\[
\text{freeIn (LETX (LETSTAR, b::bs, e)) y = freeIn (LETX (LET, [b], LETX (LETSTAR, bs, e))) y}
\]

• It’s easier to write freeIn if you use nested functions. Mostly the variable \( y \) doesn’t change, so you needn’t pass it everywhere. You’ll see the same technique used in the eval and ev functions in the chapter, as well as the model solution for eval on the continuations homework.

• If you can apply what you have learned on the scheme and hofs assignments, you should be able to write improve on one line, without using any explicit recursion.

• Let the compiler help you: compile early and often.

• Once you have the interpreter working, test it by running it on your solutions from the scheme, hofs, or continuation assignments. (It’s possible to get freeIn right but to break everything by making a bad mistake in improve. Testing on a full solution set will reveal the problem.)

My implementation of freeIn is 21 lines of ML.

**Extra credit**

There are three extra-credit problems: **MULTIPLY**, **FIVES**, and **VARARGS**.

**MULTIPLY. Multiplication of natural numbers.**

Define a function

\[
\text{val mulNats : nat * nat -> nat}
\]

that multiplies two natural numbers. Multiplication obeys these algebraic laws:

\[
0 \ast n = 0 \\
n \ast 0 = 0 \\
(10 \ast m1 + d1) \ast (10 \ast m2 + d2) =
\]

\[
d1 \ast d2 + \\
10 \ast (m1 \ast d2 + m2 \ast d1) + \\
100 \ast (m1 \ast m2)
\]

Each of the summands has to be represented as a natural number:

• Number \( d1 \ast d2 \) can be computed using machine multiplication and nat0fInt.

• You can multiply \( m1 \ast d2 \) and \( m2 \ast d1 \) using nat0fInt and mulNats. The recursive call is guaranteed to terminate because at least one argument is getting smaller.

• You can multiply \( m1 \ast m2 \) using mulNats. The recursive call is guaranteed to terminate because both arguments are getting smaller.

My implementation of mulNats is seven lines of ML, plus a couple of one-line helper functions.

**FIVES. Expressions with arithmetic operators and five 5’s**

Consider the class of well-formed arithmetic computations using the numeral 5. These are expressions
formed by taking the integer literal 5, the four arithmetic operators +, -, *, and /, and properly placed parentheses. Such expressions correspond to binary trees in which the internal nodes are operators and every leaf is a 5. If you enumerate all such expressions, you can answer these questions:

- What is the smallest positive integer that cannot be computed by an expression involving exactly five 5’s?
- What is the largest prime number that can be computed by an expression involving exactly five 5’s?
- Exhibit an expression that evaluates to that prime number.

Write an ML function `reachable` of type

\[ (\text{array of \ 'a * 'a -> order}) \times (\text{array of \ 'a * 'a -> 'a}) \rightarrow \text{function of \ 'a -> \ 'a set} \]

such that `reachable (Int.compare, [op +, op -, op *, op div]) 5 5` computes the set of all integers computable using the given operators and exactly five 5’s. (You don’t have to bother giving the answers to the questions above; the first two are easily gotten with `setFold`.) My solution is under 20 lines of code. Such brevity is possible only because I rely heavily on the `setFold`, `nullset`, `addelt`, and `pairfoldrEq` functions defined earlier.

**Hints:**

- In order to be able to use `Int.compare` interactively, you will either have to run `mosml -P full` or else tell Moscow ML interactively to load “Int”;
- Begin your function definition this way:
  ```ml
  fun reachable (cmp, operators) five n =
  (* produce set of expressions reachable with exactly n fives *)
  ```
- Use dynamic programming\(^{28}\).
- Create a list of length \(k\)-1 in which element \(i\) is a set containing all the integers that can be computed using exactly \(i\) elements. Now compute the \(k\)th element of the list by combining 1 with \(k\)-1, 2 with \(k\)-2, etcetera.
- Try doing the above by passing a list and its reverse, then use `pairfoldrEq` with a suitable function.
- The initial list contains a set with exactly one element (in the example above, 5).
- Make sure your solution has the completely general type given above, so you could use it with different operations and with different representations of numbers.

**VARARGS. Variadic functions in Scheme.**

Extend μScheme to support procedures with a variable number of arguments. This is Exercise 8 on page 391 of *Build, Prove, and Compare*.

### Avoid common mistakes

It’s a common mistake to use any of the functions `length`, `hd`, and `tl`. Instant **No Credit**.

\(^{28}\) [readings/dynamic.html](../readings/dynamic.html)
If you redefine a type that is already in the initial basis, code will fail in baffling ways. (If you find yourself baffled, exit the interpreter and restart it.) If you redefine a function at the top-level loop, this is fine, unless that function captures one of your own functions in its closure.

Example:

```ml
fun f x = ... stuff that is broken ...
fun g (y, z) = ... stuff that uses ‘f’ ...
fun f x = ... new, correct version of ‘f’ ...
```

You now have a situation where **g is broken, and the resulting error is very hard to detect**. Stay out of this situation; instead, load fresh definitions from a file using the use function.

**Never put a semicolon after a definition.** I don’t care if Jeff Ullman does it—don’t you do it. It’s wrong! Write a semicolon only when you are deliberately using imperative features.

It’s a common mistake to become very confused by **not knowing where you need to use op**. Ullman covers op in Section 5.4.4, page 165.

It’s a common mistake to **include redundant parentheses in your code**. To avoid this mistake, use the checklist in the section Expressions VIII (Parentheses) in *Learning Standard ML*.29

It’s a common mistake to do both your pair work and your solo work in the same directory. The submit scripts will balk.

It’s not a common mistake, but it can be devastating: when you’re writing a type variable, be sure to use an ASCII quote mark, as in ’a, not with a Unicode right quote mark, as in ‘a. As you may remember from class, some text editors, web browsers, or Bluetooth keyboards may use or display Unicode without being asked. Thanks, Apple!

It’s not a common mistake, but do not copy Unit.sml into your submission directory—you won’t be able to submit.

**What to submit and how to submit it**

There is no README file for this assignment.

**Submitting your individual work**

For your individual work, please submit the files cqs.ml.txt and warmup.sml. If you have implemented mulNats, please include it in warmup.sml. If you have done either of other the extra-credit problems, please submit them as varargs.sml or fives.sml.

In comments at the top of your warmup.sml file, please include your name and the names of any collaborators, and a note about any extra-credit work you have done.

As soon as you have a warmup.sml file, run submit105-ml-solo to submit a preliminary version of your work. As you edit your files, keep submitting; we grade only the last submission.

---

29 ../readings/ml.html#expressions-viii-parentheses
Submitting your improved μScheme interpreter

For your improved μScheme interpreter, which you may have done with a partner, please submit the file mlscheme-improved.sml, using the script submit105-ml-pair.

How your work will be evaluated

The criteria are mostly the same as for the scheme and hofs assignments, but because the language is different, we’ll be looking for indentation and layout as described in the Style Guide for Standard ML Programmers\(^{30}\).

\(^{30}\)../handouts/mlstyle.pdf