

`(power x 0) == 1`
`(power x (+ n 1)) == (* (power x n) x)`

`(power x 0) == 1`
`(power x (+ (* 2 m) 0)) == (square (power x m))`
`(power x (+ (* 2 m) 1)) == (* x (square (power x m)))`

`(all-fours? d) == (= d 4), [d is digit]`
`(all-fours? (+ (* 10 m) d)) ==`
`(and (= d 4) (all-fours? m)), where m != 0`

`(has-digit? d d) == 1`
`(has-digit? d d') == 0, where d differs from d'`
`(has-digit? (+ (* 10 m) d) d') ==`
`(or (= d d') (has-digit? m d')), where m != 0`

Bloom's taxonomy (Metacognition 10)

Cognitive actions:

- 1. Remember**
- 2. Understand**
- 3. Apply**
- 4. Analyze**
- 5. Evaluate**
- 6. Create**

Operational semantics

Cognitive actions:

1. Remember
2. Understand
3. Apply
4. **Analyze**
5. Evaluate
6. Create

Concrete syntax for Impcore

Definitions and expressions:

```
def ::= (define f (x1 ... xn) exp)    ;; "true" defs
      | (val x exp)
      | exp
      | (use filename)                ;; "extended" defs
      | (check-expect exp1 exp2)
      | (check-assert exp)
      | (check-error exp)
```

```
exp ::= integer-literal
      | variable-name
      | (set x exp)
      | (if exp1 exp2 exp3)
      | (while exp1 exp2)
      | (begin exp1 ... expn)
      | (function-name exp1 ... expn)
```

How to define behaviors inductively

Expressions only

Base cases (plural): numerals, names

Inductive steps: compound forms

- **To determine behavior of a compound form, look at behaviors of its parts**

First, simplify the task of specification

What's different? What's the same?

```
x = 3;
```

```
(set x 3)
```

```
while (i * i < n)
```

```
(while (< (* i i) n)
```

```
  i = i + 1;
```

```
  (set i (+ i 1)))
```

Abstract away gratuitous differences

(See the bones beneath the flesh)

Abstract syntax

Same inductive structure as BNF grammar
(related to proof system)

More uniform notation

Good representation in computer

Concrete syntax: sequence of symbols

Abstract syntax: ???

The abstraction is a tree

The abstract-syntax tree (AST):

```
Exp = LITERAL (Value)
     | VAR      (Name)
     | SET      (Name name, Exp exp)
     | IFX      (Exp cond, Exp true, Exp false)
     | WHILEX   (Exp cond, Exp exp)
     | BEGIN    (Explist)
     | APPLY    (Name name, Explist actuals)
```

One kind of “application” for both user-defined and primitive functions.

In C, trees are fiddly

```
typedef struct Exp *Exp;
typedef enum {
    LITERAL, VAR, SET, IFX, WHILEX, BEGIN, APPLY
} Expalt;          /* which alternative is it? */

struct Exp { // only two fields: 'alt' and 'u'!
    Expalt alt;
    union {
        Value literal;
        Name var;
        struct { Name name; Exp exp; } set;
        struct { Exp cond; Exp true; Exp false; } ifx;
        struct { Exp cond; Exp exp; } whilex;
        Explist begin;
        struct { Name name; Explist actuals; } apply;
    } u;
};
```

Let's picture some trees

An expression:

```
(f x (* y 3))
```

(Representation uses `ExprList`)

A definition:

```
(define abs (n)
  (if (< n 0) (- 0 n) n))
```

Behaviors of ASTs, part I: Atomic forms

Numeral: stands for a **value**

Name: stands for what?

In Impcore, a name stands for a value

Environment associates each **variable** with one **value**

Written $\rho = \{x_1 \mapsto n_1, \dots, x_k \mapsto n_k\}$,
associates variable x_i with value n_i .

Environment is **finite map**, aka **partial function**

$x \in \text{dom } \rho$ x is defined in environment ρ

$\rho(x)$ the value of x in environment ρ

$\rho\{x \mapsto v\}$ extends/modifies environment ρ to map x to v

Environments in C, abstractly

An abstract type:

```
typedef struct Valenv *Valenv;  
  
Valenv mkValenv(Namelist vars, Valuelist vals);  
bool isvalbound(Name name, Valenv env);  
Value fetchval (Name name, Valenv env);  
void bindval   (Name name, Value val, Valenv env);
```

“Environment” is pointy-headed theory

You may also hear:

- **Symbol table**
- **Name space**

Influence of environment is “scope rules”

- **In what part of code does environment govern?**

Find behavior using environment

Recall

`(* y 3) ; ;` what does it mean?

Your thoughts?

Impcore uses three environments

Global variables ξ

Functions ϕ

Formal parameters ρ

There are **no local variables**

- Just like `awk`; if you need temps, use extra formal parameters
- For homework, you'll add local variables

Function environment ϕ not shared with variables—just like Perl