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def ::= (val variable-name exp)
      | exp
      | (define function-name (formals) exp)
      * | (record record-name [ {field-name} ])
      | (use file-name)
      | unit-test

unit-test ::= (check-expect exp exp)
              | (check-assert exp)
              | (check-error exp)

exp ::= literal
      | variable-name
      | (set variable-name exp)
      | (if exp exp exp)
      | (while exp exp)
      | (begin {exp})
      | (exp {exp})
      | (let-keyword ( { [ variable-name exp ] } ) exp)
      | (lambda (formals) exp)
      * | (&& {exp}) | (|| {exp})
      * | (cond { [ question-exp answer-exp ] })
      * | (when exp {exp}) | (unless exp {exp})

let-keyword ::= let | let* | letrec

formals ::= { variable-name }

literal ::= numeral | #t | #f | 'S-exp | (quote S-exp)

S-exp ::= symbol-name | numeral | #t | #f | ( { S-exp } )

numeral ::= token composed only of digits, possibly prefixed with a plus
           or minus sign

*-name ::= token that is not a bracket, a numeral, or one of the “re-
          served” words shown in typewriter font
  
```

Tokens are as in Impcore, except that if a quote mark ' occurs at the beginning of a token, it is a token all by itself; e.g., 'yellow is two tokens.

Each quoted S-expression is converted to a literal value by the parser. And each record definition is expanded to a sequence of true definitions, also by the parser; in other words, a record definition is syntactic sugar (Section 1.8 on page 68), as marked by the *. Five forms of conditional expression are also syntactic sugar. All the other forms are handled by the eval function.

Figure 2.2: Concrete syntax of μ Scheme